

# Performance Bounds for Erasure, List and Decision Feedback Schemes with Linear Block Codes

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## Abstract

A message independence property and some new performance upper bounds are derived in this work for erasure, list and decision-feedback schemes with linear block codes transmitted over memoryless symmetric channels. Similar to the classical work of Forney, this work is focused on the derivation of some Gallager-type bounds on the achievable tradeoffs for these coding schemes, where the main novelty is the suitability of the bounds for both random and structured linear block codes (or code ensembles). The bounds are applicable to finite-length codes and to the asymptotic case of infinite block length, and they are applied to low-density parity-check code ensembles.

## Index Terms

Automatic repeat request (ARQ), erasures, error exponents, decision feedback, linear codes, list decoding, low-density parity-check (LDPC) codes.

## I. INTRODUCTION

Exponential error bounds were derived and studied by Forney [15], referring to the following two situations:

- 1) A decoder is allowed not to make a decision on a received signal, or rejecting all estimates; this output is called an *erasure*. When a decision is made, the event where the decoder decision is incorrect is called an *undetected error*.
- 2) A decoder is allowed to make more than one estimate of the received signal. The output of this decoder forms a list of codewords, and the event where the transmitted message is not on the list is called a *list error event*.

Throughout this paper, decoding rules for these two situations are called *generalized decoding rules* since they apply to the general setting where the decoder does not necessarily need to make a single decision about the codeword which was sent. As explained in [15], erasure and list options may be useful when the transmitted data contains some redundancy, when a feedback channel is available, or when several stages of coding (e.g., concatenation) are used. The size of the decoded list in [15] is allowed to vary according to the received signal. This decoding rule differs from [14] and [36] where the size of the list is predetermined and fixed.

By allowing a decoder to increase the probability of erasures in the first case, the undetected error probability can be reduced. In the second case, by increasing the decoder list, the list error probability can be reduced. The optimum decoding rules with respect to these tradeoffs were provided in [15] and they were analyzed via the derivation of exponential bounds for random codes. Some sub-optimal decoding rules are analyzed in [20] and [21] for random codes via a similar bounding technique, and the random coding error exponents under optimal and

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sub-optimal decoding rules are compared. It is noted that the considered decoding rules are studied with respect to a given code, and finding the optimal codes for these scenarios remains an open problem.

Much of the current literature on the performance analysis of codes is focused on maximum-likelihood (ML) decoding (see, e.g., [26] and references therein). Lower bounds on the error exponents for fully-random block codes under generalized decoding rules are derived in [2], [15], [25], and [33]. Achievable error exponents are provided in [30] and [31] for random codes with constant composition under some suboptimal decoding rules. An upper bound on the error exponent under fixed-size list-decoding is provided in [28]. The error performance under fixed-size list-decoding is studied for specific codes in [1], [4] and [24] where the communication is assumed to take place over an AWGN channel. Additional (suboptimal) decoding rules with erasures are analyzed in [10] and [11].

The analysis of error probabilities under generalized decoding rules with erasures, enables the study of coded communications with a noiseless decision feedback. Specifically, it is assumed that decoding erasures are followed by a repeat request over a noiseless and immediate feedback channel. Such schemes are often referred to as hybrid automatic repeat request (ARQ) systems. Unlike channel capacity for discrete memoryless channels, which is not affected by feedback (see for example [7, p. 216]), a significant improvement is demonstrated in [15] for the achievable error exponents. In this respect, the reader is also referred to [19] where the error exponents of hybrid ARQ schemes with limited retransmissions are studied. The effect of feedback was also considered in [6], and it was shown to significantly reduce the block error probability for discrete memoryless channels.

In this paper, upper bounds on the error probabilities under generalized decoding rules are provided for linear block codes over memoryless symmetric channels. Both optimal and suboptimal decoding rules are considered. When variable-size list-decoding is considered, upper bounds on the expected size of the decoded list and the associated error probability under list decoding are jointly derived. In addition, upper bounds on the list error probability are introduced for linear block codes when the size of the list is fixed. The bounds derived in this work are applicable to the performance analysis of specific codes, and code ensembles, via their (average) distance spectra. The bounds are suitable for finite block lengths and also for the asymptotic case of an infinite block length. The provided results are exemplified for two coding schemes: Fully-random linear block codes, and regular, binary and non-binary, low-density parity-check (LDPC) code ensembles with finite block lengths. An application to coded hybrid-ARQ schemes is also exemplified.

This paper is structured as follows: The definitions of channel symmetry, generalized decoding rules, and some of their basic properties, are provided in Section II. New upper bounds under the generalized decoding rules in [15] are derived in Section III. Error performance of suboptimal decoding rules are provided in Sections IV and V. Some technical details are relegated to the appendices.

## II. CHANNEL SYMMETRY, GENERALIZED DECODING, AND MESSAGE INDEPENDENCE

In this section we introduce some definitions, examples, and statements related to channel symmetry, Forney's generalized decoding rule [15], and sub-optimal versions ([2] and [15]), as well as list decoding rules ([14] and [36]). A message independence property is stated for these decoding rules, which is used for the simplification of the analysis.

Let  $\mathcal{X} = \{x_0, x_1, \dots, x_{q-1}\}$  be a given alphabet with a cardinality  $q$ . We assume an addition operation (+) over the alphabet  $\mathcal{X}$  for which  $\{\mathcal{X}, +\}$  forms an Abelian group. Let  $x_0 = 0$  be the additive identity of this group. In addition, let  $\mathcal{Y}$  be a given discrete (or continuous) alphabet. We assume a memoryless channel, and denote the channel transition probability (or probability density, respectively) function by  $p(y|x)$ , where  $x \in \mathcal{X}$  and  $y \in \mathcal{Y}$ .

**Definition 1 (Channel symmetry).** A memoryless channel which is characterized by a transition probability  $p$ , an input-alphabet  $\mathcal{X}$  and a discrete output alphabet  $\mathcal{Y}$  is *symmetric* if there exists a function  $\mathcal{T} : \mathcal{Y} \times \mathcal{X} \rightarrow \mathcal{Y}$  which satisfies the following properties:

- 1) For every  $x \in \mathcal{X}$ , the function  $\mathcal{T}(\cdot, x) : \mathcal{Y} \rightarrow \mathcal{Y}$  is bijective.
- 2) For every  $x_1, x_2 \in \mathcal{X}$  and  $y \in \mathcal{Y}$ , the following equality holds:

$$p(y|x_1) = p(\mathcal{T}(y, x_2 - x_1)|x_2). \quad (1)$$

**Remark 1.** For channels whose output alphabet is continuous, an additional requirement on the mapping  $\mathcal{T}$  is that its Jacobian is equal to 1.<sup>1</sup> In this case, the condition in (1) implies that

$$\int p(y|x_1) dy = \int p(\mathcal{T}(y, x_2 - x_1)|x_2) dy.$$

**Example 1 (Memoryless binary-input output symmetric channels).** Consider a memoryless binary-input output-symmetric (MBIOS) channel. Setting

$$\mathcal{T}(y, x) = \begin{cases} y & x = 0 \\ -y & x = 1 \end{cases}$$

then Definition 1 coincides with the standard definition of MBIOS channels.

Let  $\mathcal{C} = \{\mathbf{x}_m\}_{m=1}^{q^k}$  be a linear block code whose generator matrix is a  $k \times n$  full-rank matrix with entries over  $\mathcal{X}$ . The decoding rules studied in this paper are specified in terms of decision regions  $\Lambda_m$ ,  $1 \leq m \leq q^k$ , which are all subsets of  $\mathcal{Y}^n$ . The conditional error probability of the  $m$ -th message is given by

$$P_{e|m} = \sum_{\mathbf{y} \in \Lambda_m^c} p(\mathbf{y}|\mathbf{x}_m) \quad (2)$$

where  $\Lambda_m$  forms the decision region for the  $m$ -th codeword, and the superscript ‘c’ stands for the complementary set. The decision region of the  $m$ -th codeword under ML decoding gets the form

$$\Lambda_m = \{\mathbf{y} : p(\mathbf{y}|\mathbf{x}_m) > p(\mathbf{y}|\mathbf{x}_{m'}), \forall m' \neq m\} \quad (3)$$

where ties are resolved randomly with equal probability. Assuming equal a-priori probabilities for the transmitted messages, the ML decoding rule minimizes the error probability given in (2). A well-known result for binary linear block codes operating over MBIOS channels is that their error probability under ML decoding is independent of the transmitted codeword. This enables a great simplification in the analysis by assuming that the all-zero codeword is transmitted. This result is generalized in [22] for non-binary linear block codes whose transmission takes place over memoryless symmetric channels with discrete input alphabet.

When generalized decoding rules are considered, the decision regions  $\Lambda_m$  are not necessarily disjoint nor they include all the possible received vectors. The former case corresponds to decoding rules with a possibly *variable* list-size, and the latter case corresponds to decoding with erasures. A list is produced by the decoder where the received vector may possibly belong to more than one decision region. An erasure event is declared by the decoder when the received vector does not belong to any decision region. These concepts were first introduced in [15]. When generalized decoding rules are allowed, the conditional block error probability  $P_{e|m}$  in (2) stands for the probability of either an undetected error or an erasure. When the decision regions are disjoint, the conditional undetected error probability is given by

$$P_{ue|m} = \sum_{m' \neq m} \sum_{\mathbf{y} \in \Lambda_{m'}} p(\mathbf{y}|\mathbf{x}_m). \quad (4)$$

In addition, let  $P_{x|m}$  denote the conditional probability of an erasure event given that  $\mathbf{x}_m$  is transmitted. Then

$$P_{x|m} = P_{e|m} - P_{ue|m}.$$

In the case where list decoding is considered, the decision regions are not disjoint, and  $P_{ue|m}$  as given in (4) is no longer a probability. However the RHS of (4) equals the conditional expectation of the number of incorrect codewords in the list (the same notation,  $P_{ue|m}$ , is used in both cases to simplify the statement of the following results). The optimum decoding rule with respect to the tradeoff between the error and the undetected error event is derived in [15].

<sup>1</sup>It is possible to use a generalized definition for both discrete and continuous output alphabets using the notion of unitary functions, as done for example in [35, Section III-A].

**Definition 2 (Forney's generalized decoding).** Consider a block code over an alphabet  $\mathcal{X}$ , and let  $\{\mathbf{x}_m\}$  denote its codebook. The generalized decoding rule is defined by the following decision regions:

$$\Lambda_m = \left\{ \mathbf{y} \in \mathcal{Y}^n : \frac{\Pr(\mathbf{y}, \mathbf{x}_m)}{\sum_{m' \neq m} \Pr(\mathbf{y}, \mathbf{x}_{m'})} \geq e^{nT} \right\} \quad (5)$$

where  $m$  is the index of the codeword,  $T \in \mathbb{R}$  is a parameter,  $\Pr(\mathbf{y}, \mathbf{x}_m)$  denotes the joint probability that  $\mathbf{x}_m$  is the transmitted codeword and  $\mathbf{y}$  is the received vector, and the summation is over all codewords except for  $\mathbf{x}_m$ .

**Remark 2.** The decision region in (5) can be expressed equivalently in the form

$$\Lambda_m = \left\{ \mathbf{y} \in \mathcal{Y}^n : \Pr(\mathbf{x}_m | \mathbf{y}) \geq \frac{e^{nT}}{1 + e^{nT}} \right\}. \quad (6)$$

Note that for  $T = 0$ , this decision region includes all the vectors  $\mathbf{y} \in \mathcal{Y}^n$  for which  $\Pr(\mathbf{x}_m | \mathbf{y}) \geq \frac{1}{2}$ . The a-posteriori probability of  $\mathbf{x}_m$ , given that  $\mathbf{y} \in \Lambda_m$  is received, is therefore larger than the a-posteriori probability for any other codeword. Hence, if a codeword is selected according to the decoder with the decision regions in (6) with  $T = 0$ , then the same decision is made by a MAP decoder (as no other codeword can get an a-posteriori probability larger than  $\frac{1}{2}$ ). This implies that the undetected error exponent for the decoder in (6) with  $T = 0$  cannot be smaller than the error exponent of an ML decoder with equally-likely codewords. Interestingly, as will be shown later, we get the same lower bound on the error exponents for both decoders. Moreover, it is shown that for  $T = 0$  the bounds for the undetected error event and erasures coincides.

**Remark 3.** The threshold parameter  $T$  in (5) controls the tradeoff between erasures and undetected errors (or average list size and decoding error). Setting  $T > 0$  guarantees that the decision regions  $\Lambda_m$  are disjoint.

**Proposition 1 (Forney's generalized decoding [15]).** Assume that the decoding of a block code is carried according to the generalized decoding rule in Definition 2. Then, there is no other decoding rule that simultaneously gives a lower error probability and a lower undetected error probability (or an average number of incorrect codewords when list decoding is considered).

**Remark 4 (On optimal generalized decoding of convolutional codes).** Optimal generalized decoding of convolutional codes, whose transmission takes place over memoryless channels, is provided in [23]. This algorithm is based on the decision regions in (5). Specifically, the algorithm is based on a modification of the standard Viterbi algorithm, where the denominator in (5) is evaluated recursively. The optimality of the algorithm in [23] is based on the optimality in Proposition 1.

The following proposition generalizes the message independence property for the case of generalized decoding:

**Proposition 2 (Message independence property for optimal generalized decoding).** Let  $\mathcal{C}$  be a linear block code whose transmission takes place over a memoryless and symmetric channel. Then, the block error probability and the undetected error probability, under the generalized decoding rule in Definition 2, are independent of the transmitted codeword.

*Proof:* See Appendix A. ■

**Remark 5.** In the case where list decoding is considered (i.e., the decision regions are not disjoint), then Proposition 2 holds when we refer to the conditional expectation of the number of incorrect messages in the list produced by the generalized decoding rule, instead of the undetected error probability.

The following suboptimal decoding rule is suggested in [15] for the case of decoding with erasures:

**Definition 3 (Likelihood Ratio (LR) Decoding).** Consider a block code over the alphabet  $\mathcal{X}$ , and let  $\{\mathbf{x}_m\}$  denote its codebook. The LR decoding rule is defined by the following decision regions:

$$\Lambda_m^{\text{LR}} = \left\{ \mathbf{y} \in \mathcal{Y}^n : \frac{\Pr(\mathbf{y}, \mathbf{x}_m)}{\Pr(\mathbf{y}, \mathbf{x}_{m_2})} \geq e^{nT} \right\} \quad (7)$$

where  $m$  is a codeword index,  $T > 0$  is a parameter,  $\Pr(\mathbf{y}, \mathbf{x}_m)$  denotes the joint probability that  $\mathbf{x}_m$  is the transmitted codeword and  $\mathbf{y}$  is the received vector, and  $m_2 = m_2(\mathbf{y})$  denotes the second most probable codeword for each received vector  $\mathbf{y}$ .

**Remark 6.** It is observed in [15] that the LR decoding rule may be a good approximation to the optimal regions in (5), since the second most likely codeword is usually much more probable than the rest of the codewords (excluding the most probable codeword). It is also noted in [15] that this suboptimal decoding rule is of practical utility.

**Example 2 (Suboptimal generalized decoding).** Consider the transmission of a binary linear block code over a BSC. Given a received vector  $\mathbf{y} \in \{0, 1\}^n$ , the decoded codeword is  $\mathbf{x}$  if and only if

$$d_H(\mathbf{x}', \mathbf{y}) - d_H(\mathbf{x}, \mathbf{y}) > 2\tau n \quad (8)$$

for all codewords  $\mathbf{x}' \neq \mathbf{x}$ , where  $d_H(\mathbf{x}, \mathbf{y})$  denoted the Hamming distance between  $\mathbf{x}$ , and  $\mathbf{y}$ , and  $\tau \geq 0$  is an arbitrary parameter. Otherwise, an erasure is declared. It is easily verified that this rule is a particular case of (7). The error exponents for this setting are studied in [2].

The following proposition obtains a message independence property for the suboptimal decoding rule in Definition 3:

**Proposition 3 (Message independence property for (suboptimal) LR decoding).** Let  $\mathcal{C}$  be a linear block code whose transmission takes place over a memoryless and symmetric channel. Then, the block error probability and the undetected error probability, under the suboptimal decoding rule in (7), are independent of the transmitted codeword.

*Proof:* See Appendix B. ■

The following definition considers list decoding with a fixed size. Such a decoding rule is based on a fixed size of the list (instead of a variable list size which characterizes the decoding rule in Definition 2 with  $T < 0$ ).

**Definition 4 (Fixed-size list-decoding).** Consider a block code over an alphabet  $\mathcal{X}$ , and let  $\{\mathbf{x}_m\}$  denote its codebook. Given a fixed list size  $L$ , the list-decoder is a mapping from the set of all possible received vectors  $\mathcal{Y}^n$  to the set of all possible lists of  $L$  codewords. This mapping produces the list whose likelihoods are the highest among all other codewords. That is, given a received vector  $\mathbf{y}$ , a codeword  $\mathbf{x}_m$  is in the list if  $p(\mathbf{y}|\mathbf{x}_m) > p(\mathbf{y}|\mathbf{x}_{m'})$  for all  $m' \neq m$  except for at most  $L - 1$  other possible codewords.

Assuming that the codeword  $\mathbf{x}_m$  is transmitted, a block error event is occurred by the fixed-size list-decoding rule in Definition 4, if the list produced by the decoder does not include the transmitted codeword  $\mathbf{x}_m$ . The following proposition is analogous to the message independence property in Propositions 2 and 3:

**Proposition 4 (Message independence property for fixed-size list-decoding).** Let  $\mathcal{C}$  be a linear block code whose transmission takes place over a memoryless and symmetric channel. Then, the block error probability, under the fixed-size list-decoding is independent of the transmitted codeword.

*Proof:* See Appendix C. ■

### III. UPPER BOUNDS UNDER OPTIMAL GENERALIZED DECODING

The transmission of block codes (not necessarily linear) is first considered. In addition, throughout the paper, all codewords are assumed to have a uniform a-priori probability.

**Proposition 5.** Consider the transmission of a code  $\mathcal{C}$  with a block length  $n$  and  $M$  codewords, and let  $p(\mathbf{y}|\mathbf{x})$  designate the transition probability of the channel where  $\mathbf{x} \in \mathcal{C}$  is the transmitted codeword and  $\mathbf{y} \in \mathcal{Y}^n$  is the received vector. Then, the conditional block error probability ( $P_{e|m}$ ) and the average undetected error probability ( $P_{ue}$ ) under the generalized decoding rule in (5) satisfy

$$P_{e|m} \leq e^{nsT} D_B(m, G_n^m, s, \rho) \quad (9)$$

$$P_{ue} \leq e^{n(s-1)T} \frac{1}{M} \sum_{m=1}^M D_B(m, G_n^m, s, \rho) \quad (10)$$

where  $0 \leq s \leq \rho \leq 1$  are real-valued parameters,  $G_n^m$  is an arbitrary non-negative function over  $\mathcal{Y}^n$  which possibly depends on the codeword  $\mathbf{x}_m$ ,  $1 \leq m \leq M$ , and

$$D_B(m, G_n^m, s, \rho) \triangleq \left( \sum_{\mathbf{y}} G_n^m(\mathbf{y}) p(\mathbf{y}|\mathbf{x}_m) \right)^{1-\rho} \left( \sum_{m' \neq m} \sum_{\mathbf{y}} p(\mathbf{y}|\mathbf{x}_{m'}) G_n^{m'}(\mathbf{y})^{1-\frac{1}{\rho}} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^{\rho}. \quad (11)$$

*Proof:* See Appendix D. ■

**Remark 7.** Bounds (9) and (10) in Proposition 5 may be considered as a generalization of the DS2 bound ([9], [27], [26]). In fact, setting  $T = 0$  in (9) reproduces the DS2 bound under ML decoding. Note however that for  $T = 0$ , the decision regions in (5) do not coincide with those under ML decoding (e.g., in the former case there are erasures).

The following corollary is a particularization of Proposition 5 for fully random block codes whose transmission takes place over memoryless channels. The corollary reproduces the exponential upper bounds as in [15, Th. 2].

**Corollary 1 (Random coding error exponents under optimum generalized decoding).** Consider the transmission of block codes over a memoryless communication channel with a transition probability law  $p$ . Then, under the notation in Proposition 5, there exists a block code which simultaneously satisfies

$$P_e \leq e^{-nE_1(R,T)} \quad (12)$$

$$P_{ue} \leq e^{-nE_2(R,T)} \quad (13)$$

where  $R = \ln M/n$  is the code rate (in nats per channel use),

$$E_1(R, T) \triangleq \max_{0 \leq s \leq \rho \leq 1, q_X} \left( E_0(s, \rho, q_X) - \rho R - sT \right) \quad (14)$$

$$E_2(R, T) \triangleq E_1(R, T) + T$$

$$E_0(s, \rho, q_X) \triangleq -\ln \sum_{y \in \mathcal{Y}} \left\{ \left( \sum_{x \in \mathcal{X}} q_X(x) p(y|x)^{1-s} \right) \left( \sum_{x \in \mathcal{X}} q_X(x) p(y|x)^{\frac{s}{\rho}} \right)^{\rho} \right\} \quad (15)$$

and  $q_X$  is a probability distribution over  $\mathcal{X}$ .

*Proof:* See Appendix E. ■

The bounds in Corollary 1 are derived in [15] without relying on tilting measures. The current derivation relies on the DS2 bound which makes use of tilting measures and Jensen's inequality. It is noted in [15] that setting  $T = 0$  in Corollary 1, provides the random coding error exponent of Gallager [17]. Hence, as is mentioned in [15], the random coding error exponent is attainable not only under ML decoding, but also under the generalized decoding rule in (5) with  $T = 0$ . The following proposition is a particularization of Proposition 5 for linear block codes.

**Proposition 6.** Consider an  $(n, k)$  linear block code  $\mathcal{C}$  whose transmission takes place over a memoryless symmetric channel. Assume that the channel input and output alphabets are  $\mathcal{X}$  and  $\mathcal{Y}$ , respectively, and let  $p$  be the transition probability of the channel. Then, the block error probability  $P_e$  and the undetected error probability  $P_{ue}$  under the generalized decoding rule in (5), satisfy

$$P_e \leq e^{nsT} D(g, s, \rho) \quad (16)$$

$$P_{ue} \leq e^{-n(1-s)T} D(g, s, \rho) \quad (17)$$

where  $g : \mathcal{Y} \rightarrow \mathbb{R}$  is an arbitrary non-negative real-valued function,  $0 \leq s \leq \rho \leq 1$  are arbitrary parameters, and

$$D(g, s, \rho) \triangleq \left( \sum_{y \in \mathcal{Y}} g(y) p(y|0) \right)^{n(1-\rho)} \left( \sum_{m' \neq 0} \prod_{i=1}^n \sum_{y \in \mathcal{Y}} g(y)^{1-\frac{1}{\rho}} p(y|0) \left( \frac{p(y|x_{m',i})}{p(y|0)} \right)^{\frac{s}{\rho}} \right)^{\rho}. \quad (18)$$

*Proof:* See Appendix F. ■

**Remark 8.** When the decision regions are not disjoint (i.e., a list decoder is considered),  $P_{\text{ue}}$  in (17) does not denote a probability but the expected number of incorrect codewords in the decoded list. The block error probability  $P_e$  in (16) refers, in this case, to the list decoding error probability.

**Remark 9.** The parameters  $s$  and  $\rho$  in Proposition 6 may be chosen separately for the bounds in (16) and (17). However, the optimized choice of the two parameters is identical in both bounds (since they only differ in the multiplicative term  $e^{-nT}$ ).

The mathematical structure of the bound provided in the following corollary is similar to the Shulman-Feder bound (SFB) in [29]. Because of this reason, this bound may be considered as a generalization of the SFB for the generalized decoding rule in (5). To simplify the notation, the corollary is provided for the case of a binary linear block code whose transmission takes place over an MBIOS channel (the generalization of the bounds to non-binary linear block codes is performed similarly to the approach in the proof of [22, Theorem 2]).

**Corollary 2.** Consider an  $(n, k)$  binary linear block code  $\mathcal{C}$  whose transmission takes place over an MBIOS channel with a transition probability law  $p$ . Then, the block error probability  $P_e$  and the undetected error probability  $P_{\text{ue}}$  under the generalized decoding rule in (5) satisfy

$$P_e \leq e^{-n \left( E(\rho, R, \mathcal{C}) - \frac{\rho T}{1+\rho} \right)} \quad (19)$$

$$P_{\text{ue}} \leq e^{-n \left( E(\rho, R, \mathcal{C}) + \frac{T}{1+\rho} \right)} \quad (20)$$

where  $0 \leq \rho \leq 1$  is an arbitrary real-valued parameter,  $R \triangleq \left(\frac{k}{n}\right) \cdot \ln 2$  is the code rate (in nats per channel use),

$$E(\rho, R, \mathcal{C}) \triangleq E_0(\rho) - \rho \left( R + \frac{\ln(\alpha(\mathcal{C}))}{n} \right) \quad (21)$$

$$E_0(\rho) \triangleq -\ln \left( \sum_y \left( \frac{1}{2} p(y|0)^{\frac{1}{1+\rho}} + \frac{1}{2} p(y|1)^{\frac{1}{1+\rho}} \right)^{1+\rho} \right) \quad (22)$$

$$\alpha(\mathcal{C}) \triangleq \max_{1 \leq i \leq n} \frac{|\mathcal{C}_i|}{2^{-(n-k)} \binom{n}{i}} \quad (23)$$

and  $|\mathcal{C}_i|$  denotes the number of codewords whose Hamming weight is  $i$ .

*Proof:* Setting  $s = \frac{\rho}{1+\rho}$ , and

$$g(y) = \left( \frac{1}{2} p(y|0)^{\frac{1}{1+\rho}} + \frac{1}{2} p(y|1)^{\frac{1}{1+\rho}} \right)^\rho p(y|0)^{-\frac{\rho}{1+\rho}} \quad (24)$$

in the bounds of Proposition 6, the proof follows in the same way as in [26, Ch. 4.4.1]. ■

**Remark 10.** In the case where the performance of an ensemble of linear block codes is of interest, repeating the derivation of Corollary 2 leads to the same upper bounds as in (19) and (20), where the cardinality  $|\mathcal{C}_i|$  in (23) is replaced with its statistical expectation over the considered ensemble, and the codebooks of this ensemble are chosen uniformly at random.

**Example 3 (Error exponents of fully random binary linear block codes).** Consider the transmission of fully random binary linear  $(n, k)$  block codes over a memoryless symmetric channel. For this particular case, the term  $\alpha(\mathcal{C})$  in (23) equals 1. As a result, it follows from Corollary 2 that the exponent of the block error probability (including erasures and undetected errors), denoted by  $E_e$ , satisfies

$$E_e \geq \max_{0 \leq \rho \leq 1} \left( E_0(\rho) - \rho R - \frac{\rho T}{1+\rho} \right) \quad (25)$$

where  $E_0(\rho)$  is defined in (22),  $R$  is the code rate (in nats per channel use), and  $T$  is the parameter of the generalized decoding rule in Definition 2. Setting  $T = 0$  in (25) reproduces the (non-expurgated) random coding error exponent

of Gallager [17]. This observation was first made by Forney for the ensemble of fully random block codes [15]. The undetected error exponent, denoted by  $E_{\text{ue}}$ , satisfies

$$E_{\text{ue}} \geq T + \max_{0 \leq \rho \leq 1} \left( E_0(\rho) - \rho R - \frac{\rho T}{1 + \rho} \right).$$

The lower bounds on the two error exponents are shown in Fig. 1 for the case of transmission over a binary-input AWGN channel with  $E_s/N_0 = -2.8$  dB (this ratio refers to the capacity limit for a rate of one-half bits per channel use). The bounds are sketched as a function of the code rate (in nats per channel use). The lower bounds on the error exponents for the case of decoding with erasures ( $T \geq 0$ ) are provided in Fig. 1(a) for  $T = 0, 0.025, 0.05, 0.1$  and  $0.15$ . For the case of decoding with a variable list-size ( $T < 0$ ), the lower bounds on the error exponents are provided in Fig. 1(b) for  $T = 0, -0.05$ , and  $-0.1$ . In addition, lower bounds on the exponent  $E_N \triangleq -(\ln N)/n$ , where  $N$  is the expected number of incorrect codewords in the decoded list, are also provided for this case. Note that the exponent  $E_N$  is negative above some rate. Only the region for which the exponent  $E_N$  is non-negative is presented; the negative part of  $E_N$ , for which an upper bound on the size of the decoded list grows exponentially with the block length, is removed.

**Definition 5 (Composition of a vector).** Let  $\mathbf{c}$  be a vector whose components are symbols in an alphabet  $\mathcal{X}$  of size  $q$ . Let us assume without loss of generality that  $\mathcal{X} = \{0, \dots, q-1\}$ . The composition of  $\mathbf{c}$ , denoted by  $\mathbf{t} = \mathbf{t}(\mathbf{c})$ , is a vector  $\mathbf{t} = (t_0, t_1, \dots, t_{q-1})$  where  $t_x$  (for  $x \in \mathcal{X}$ ) denotes the number of symbols in  $\mathbf{c}$  that are equal to  $x$ .

**Definition 6 (Complete composition spectrum).** Let  $\mathcal{C}$  be a linear block code of length  $n$  over an alphabet  $\mathcal{X}$ . The complete composition spectrum is the sequence  $\{|\mathcal{C}_{\mathbf{t}}|\}$  where  $|\mathcal{C}_{\mathbf{t}}|$  is the number of codewords whose composition is  $\mathbf{t}$ , and  $\mathbf{t}$  ranges over the set  $\mathcal{H}$  of all possible compositions over  $\mathcal{X}^n$ .

**Corollary 3.** Consider an ensemble  $\mathcal{E}$  of  $(n, k)$  linear block codes whose transmission takes place over a memoryless symmetric channel. Let  $P(l)$  denote the probability that a vector whose Hamming weight is  $l$ , forms a codeword in a randomly selected codebook from  $\mathcal{E}$ . Assume that the average composition spectrum over all the codes  $\mathcal{C}$ , uniformly selected at random from  $\mathcal{E}$  satisfies

$$\mathbb{E}[|\mathcal{C}_{\mathbf{t}}|] = P(n - t_0) \binom{n}{\mathbf{t}}. \quad (26)$$

Then, under the notation in Proposition 6, the block error probability  $P_e$  and the undetected error probability  $P_{\text{ue}}$ , satisfy

$$P_e \leq e^{\frac{n\rho T}{1+\rho}} \cdot D_s(\rho, \mathcal{C}) \quad (27)$$

$$P_{\text{ue}} \leq e^{-\frac{nT}{1+\rho}} \cdot D_s(\rho, \mathcal{C}) \quad (28)$$

where  $0 \leq \rho \leq 1$ , and

$$D_s(\rho, \mathcal{C}) \triangleq A(\rho)^{n(1-\rho)} \left( \sum_{1 \leq l \leq n} P(l) \binom{n}{l} B(\rho)^{n-l} C(\rho)^l \right)^\rho \quad (29)$$

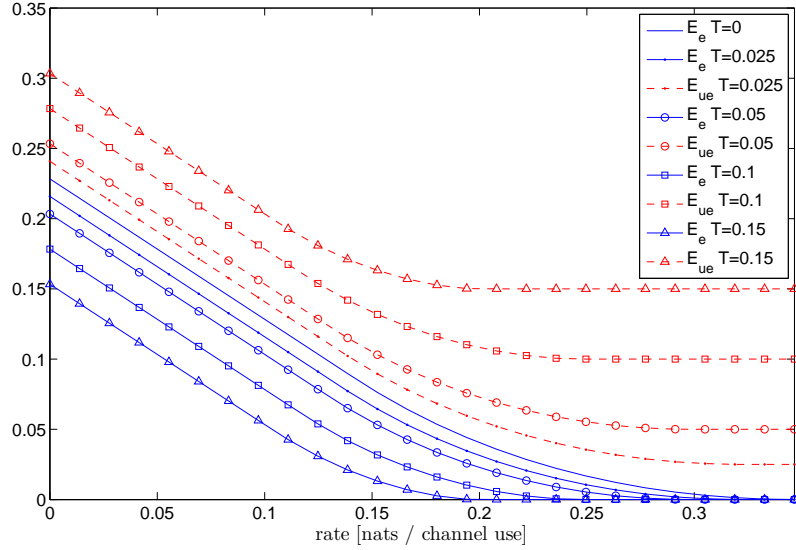
$$A(\rho) \triangleq \sum_{y \in \mathcal{Y}} \left( \frac{1}{q} \sum_{x \in \mathcal{X}} p(y|x)^{\frac{1}{1+\rho}} \right)^{1+\rho} \quad (30)$$

$$B(\rho) \triangleq \sum_{y \in \mathcal{Y}} \left( \frac{1}{q} \sum_{x \in \mathcal{X}} p(y|x)^{\frac{1}{1+\rho}} \right)^{\rho-1} \left( \frac{1}{q} \sum_{x \in \mathcal{X}} p(y|x)^{\frac{2}{1+\rho}} \right) \quad (31)$$

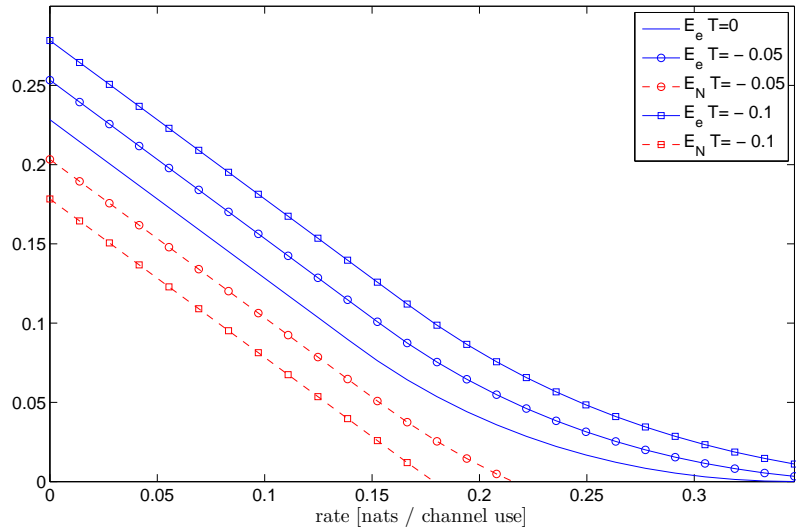
$$C(\rho) \triangleq qA(\rho) - B(\rho). \quad (32)$$

*Proof:* Setting  $s = \frac{\rho}{1+\rho}$  and choosing the tilting measure  $g$  in (24), the proof follows from Proposition 6 in the same way as in [22, Theorem 3].  $\blacksquare$





(a) Generalized decoding with erasures



(b) Generalized decoding with a variable-size list

Fig. 1: Lower bounds on the error exponents and list-size exponents for the ensemble of fully-random binary linear block codes whose transmission takes place over a binary-input AWGN channel with  $E_s/N_0 = -2.8$  dB. The lower bounds in Corollary 2 are sketched in plots (a) and (b), for the generalized decoding rule in (5) with erasures (i.e.,  $T \geq 0$ ) and with a variable list-size (i.e.,  $T < 0$ ), respectively.

**Remark 11.** For an ensemble of *binary* linear block codes, the condition in (26) is not mandatory. Repeating the derivation results in the same bounds as in Corollary 3 where the term  $P(l) \binom{n}{l}$  in (29) is replaced with the expected complete composition spectrum of the ensemble.

**Remark 12.** The bounds in Corollary 3 are tighter than those in Corollary 2. Hence, for a finite block length, the bounds in Corollary 3 are more attractive even though they lack the appealing exponential structure of the bounds in Corollary 2.

**Remark 13.** As a particular case of Remark 7, setting  $T = 0$  in (27) reproduces the upper bound on the decoding error probability of non-binary linear block codes under ML decoding in [22, Theorem 3].

The following comments concerns the numerical results shown in the examples throughout paper:

- 1) *Expurgation of codebooks*: The examples presented in this paper consider the performance of some expurgated ensembles of regular LDPC codes under generalized decoding rules. Specifically, an expurgation of the codebooks whose minimum Hamming distance is not larger than a specific value  $D_n$  is assumed. As a result, the expected complete composition spectrum  $E[|\mathcal{C}_t| | d_{\min} > D_n]$  of a codebook which is chosen uniformly at random from the expurgated ensemble, satisfies the following upper bound:

$$E[|\mathcal{C}_t| | d_{\min} > D_n] \leq \frac{E[|\mathcal{C}_t|]}{1 - \epsilon_n} \quad (33)$$

where  $E[|\mathcal{C}_t|]$  is the expected composition spectrum of the original (non-expurgated) ensemble, and

$$\sum_{t: n-t_0 \leq D_n} E[|\mathcal{C}_t|] \leq \epsilon_n. \quad (34)$$

The fraction of the removed codebooks is upper bounded by  $\epsilon_n$ . In the following examples, the value of  $\epsilon_n$  is negligible. For the (6,12) regular binary ensemble with block lengths of  $n = 504$  and  $2004$  bits,  $\epsilon_n = 3.6002 \cdot 10^{-5}$ , and  $5.5058 \cdot 10^{-8}$ , for  $D_n = 40$  and  $160$  bits, respectively. For the (8,16) regular octal alphabet ensemble with a block length of  $n = 1008$  symbols and  $D_n = 80$  symbols,  $\epsilon_n$  is around  $10^{-14}$ .

- 2) *Performance over the AWGN channel*: For the AWGN channel, the results in this paper are provided as function of the signal-to-noise ratio  $\frac{E_s}{N_0}$  where  $E_s$  is the energy per transmitted coded symbols, and  $\frac{N_0}{2}$  is the two-sided power spectral density of the additive white noise. This comment concerns both binary and non-binary codes.

**Example 4 (Error performance of binary regular LDPC code ensembles under generalized decoding with erasures).** Consider an expurgation of the binary and regular (6,12) LDPC code ensemble of Gallager [16] with a block length of  $n = 2004$  bits. In this expurgated ensemble, all the codebooks whose minimum distance is not larger than  $D_n = 160$  are removed. Upper bounds on the block error probability and the undetected error probability, under Forney's generalized decoding with erasures, are studied based on Corollary 3. The composition spectrum is upper bounded via (33) and (34), where the composition spectrum of the original (non-expurgated) regular LDPC code ensemble is evaluated using the method provided in [5], [32]). The bounds are provided for several non-negative values of  $T$  in Fig. 2, assuming that the transmission takes place over a binary-input AWGN channel. Note that if  $T = 0$ , the resulting bounds on the block error probability and the undetected error probability coincide, and they also provide an upper bound on the ML decoding error probability. The results indicate that by allowing an error probability that may be slightly higher than the upper bound on the error probability under ML decoding, significant improvement is guaranteed for the undetected error probability.

**Example 5 (Error performance of binary regular LDPC code ensembles under generalized decoding with a variable-size list).** The performance of the same expurgated ensemble as in Example 4 is studied here under Forney's generalized decoding with a variable list-size. Upper bounds on the block error probability and the expected number of incorrect codewords in the list, are evaluated based on the bounds in Corollary 3 for several non-positive values of  $T$ . These bounds are provided in Fig. 3, assuming a transmission over a binary-input AWGN channels. It is evident that only a slight improvement in the error performance is possible by using the generalized decoding rule.

Consider an expurgation of Gallager's ensemble of (8,16) regular LDPC codes [16] with an octal alphabet, and a block length of  $1008$  symbols. Consider the case where the expurgated ensemble excludes all the codebooks whose minimum distance is not larger than  $D_n = 80$ . The upper bounds on the error probabilities,

**Example 6 (Hybrid-ARQ schemes over AWGN channels with non-binary LDPC codes).** Coded communication systems with one-bit noiseless feedback are considered where a generalized decoding rule with erasures is applied at the receiver. Each decoding erasure is communicated via the feedback to the transmitter, which then retransmits its message. It is assumed that each transmitted block is decoded separately. Such a hybrid-ARQ system is described and studied in [15], where the error exponents for random coding are provided. For the case where deadlines are assumed, the error exponents for random coding are provided in [19]. For the sake of completeness, the specific details regarding the numerical results presented in this examples are surveyed in Appendix G. In this example,

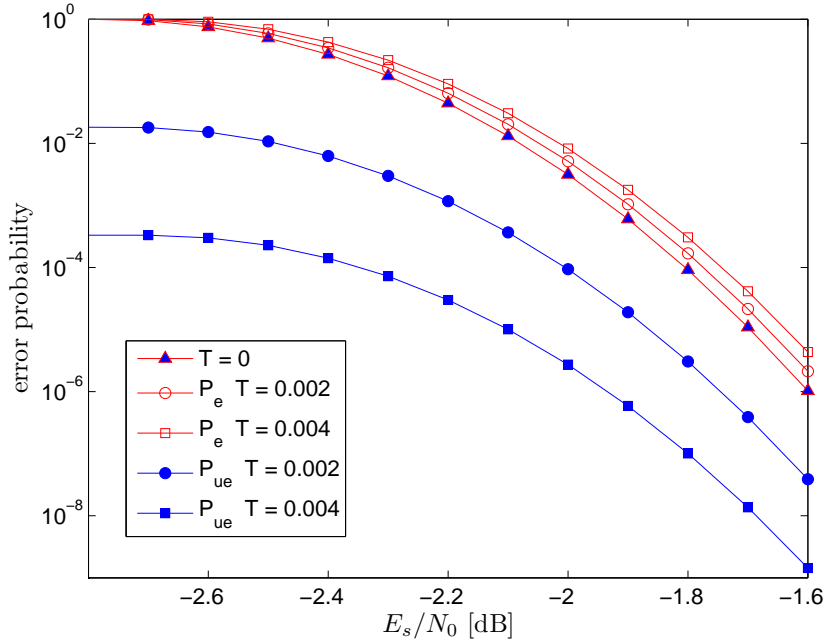


Fig. 2: Upper bounds on the block error and undetected block error probabilities under the generalized decoding rule in (5) with erasures ( $T \geq 0$ ). The transmission is assumed to take place over a binary-input AWGN channel. An expurgation of the binary and regular (6,12) LDPC code ensemble of Gallager is considered, where the block length is 2004 bits, and the parameter  $D_n$  which refers to the expurgation is set to 160 (see Example 4).

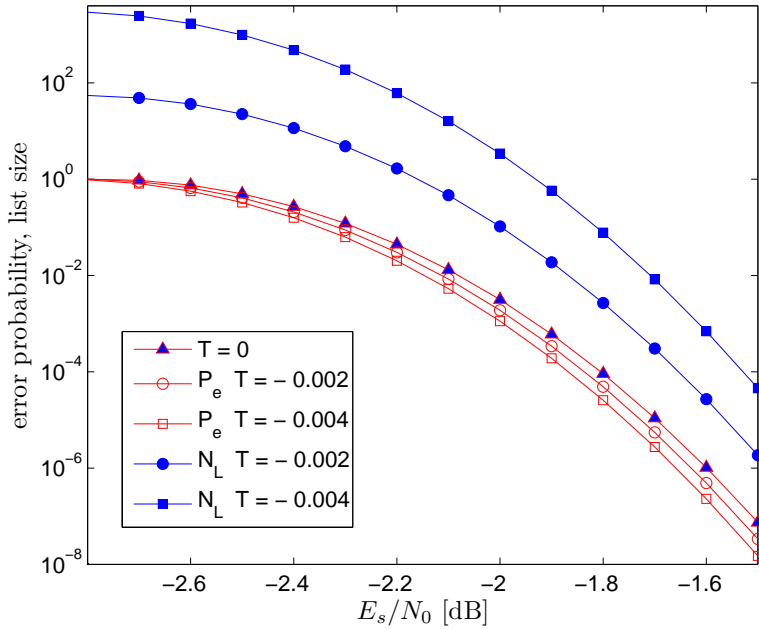


Fig. 3: Upper bounds on the block error probability and expected size of incorrect codewords in the decoded list, under the generalized decoding rule in (5) with variable-size list ( $T \leq 0$ ). The transmission is assumed to take place over a binary-input AWGN channel. An expurgation of the binary and regular (6,12) LDPC code ensemble of Gallager is considered, where the block length is 2004 bits, and the parameter  $D_n$  which refers to the expurgation is set to 160 (see Example 4).

the performance of hybrid-ARQ scheme based on octal-alphabet LDPC code ensemble is studied assuming that the transmission takes place over the AWGN channel with 8-PSK modulation. Gallager’s ensemble of (8,16) regular LDPC codes [16] with an octal alphabet, and a block length of 1008 symbols is assumed. In addition, it is assumed

that the expurgated ensemble excludes all the codebooks whose minimum distance is not larger than  $D_n = 80$ . Lower bounds on the expected rate and upper bounds on the decoding error probability are shown in Figs. 4(a) and 4(b), respectively. Schemes with and without deadlines are considered. The results show that the lower bounds on the expected rates drop considerably, below  $E_s/N_0 = 3.6$  dB. However, above this SNR, the introduction of a single-bit, noiseless and immediate feedback allows to achieve remarkable improvements in the error performance. Take for example the case where  $E_s/N_0 = 3.62$  dB where the upper bound on the error probability under ML decoding without feedback (see the curve for  $T = 0$  and  $Q = 1$ ) is around  $10^{-2}$ . For the same channel, if no deadlines are assumed, the upper bounds on the error probability are around  $2 \cdot 10^{-6}$ . When deadlines of  $Q = 2$  and 4 total retransmissions (including the first transmission) are assumed, the upper bounds on the error probability for the same channel are  $6 \cdot 10^{-4}$  and  $3 \cdot 10^{-6}$ , respectively. For all considered schemes, the expected rate deteriorates at this point by no more than 4%.

#### IV. UPPER BOUNDS UNDER SUBOPTIMAL DECODING WITH ERASURES

In this section, upper bounds on decoding error probabilities are derived for the suboptimal decoding rule in (7).

**Proposition 7.** Consider the transmission of a block code  $\mathcal{C}$  of block length  $n$  and  $M$  codewords, and let  $p(\mathbf{y}|\mathbf{x})$  designate the transition probability of the channel where  $\mathbf{x} \in \mathcal{C}$  is the transmitted codeword and  $\mathbf{y} \in \mathcal{Y}^n$  is the received vector. Then, the conditional block error probability  $P_{e|m}$ , and the conditional undetected error probability  $P_{ue|m}$ , under the suboptimal decoding rule in (7) satisfy

$$P_{e|m} \leq e^{nsT} D_B(m, G_n^m, s, \rho), \quad 0 \leq s \leq \rho \leq 1 \quad (35)$$

$$P_{ue|m} \leq e^{-nsT} D_B(m, G_n^m, s, \rho), \quad 0 \leq s \leq \rho \leq 1 \quad (36)$$

where  $D_B(m, G_n^m, s, \rho)$  is defined in (11), and  $G_n^m$  is an arbitrary non-negative function over  $\mathcal{Y}^n$  which possibly depends on the codeword  $\mathbf{x}_m$ ,  $1 \leq m \leq M$ .

*Proof:* See Appendix H. ■

**Remark 14.** The upper bound on the block error probability in (35) coincides with the upper bound on the total error probability provided in (9) under the optimal generalized decoding rule. On the other hand, the upper bounds on the undetected error probabilities under the optimal and suboptimal decoding rules in (10) and (36), respectively, are different.

The following corollary is a particularization of Proposition 7 for the ensemble of fully random block codes of length  $n$  and rate  $R$  whose transmission takes place over memoryless channels:

**Corollary 4.** Consider the transmission of block codes over a memoryless communication channel. Then, there exists a block code satisfying

$$\begin{aligned} P_e &\leq e^{-nE_1(R,T)} \\ P_{ue} &\leq e^{-nE_2^*(R,T)} \end{aligned}$$

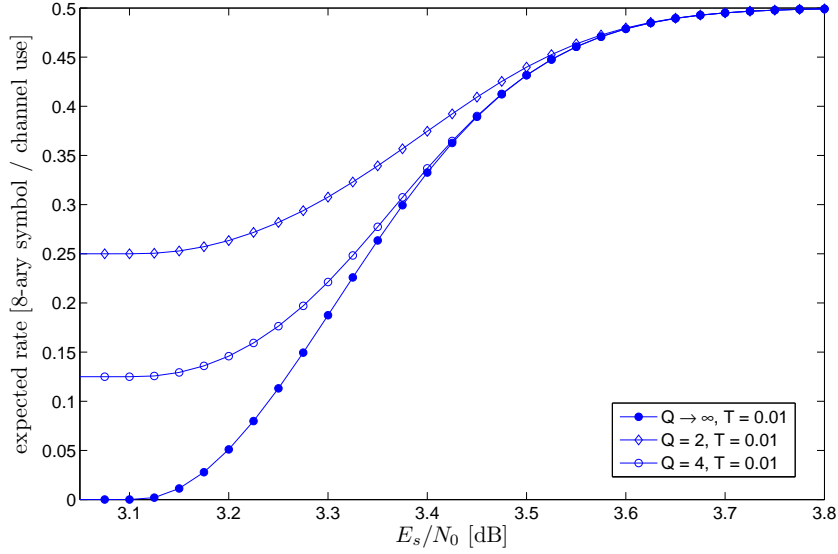
where  $R \triangleq \frac{\ln M}{n}$  is the code rate (in nats per channel use),  $E_1(R, T)$  is defined in (14),

$$E_2^*(R, T) \triangleq \max_{0 \leq s \leq \rho \leq 1, q_X} \left( E_0(s, \rho, q_X) - \rho R + sT \right)$$

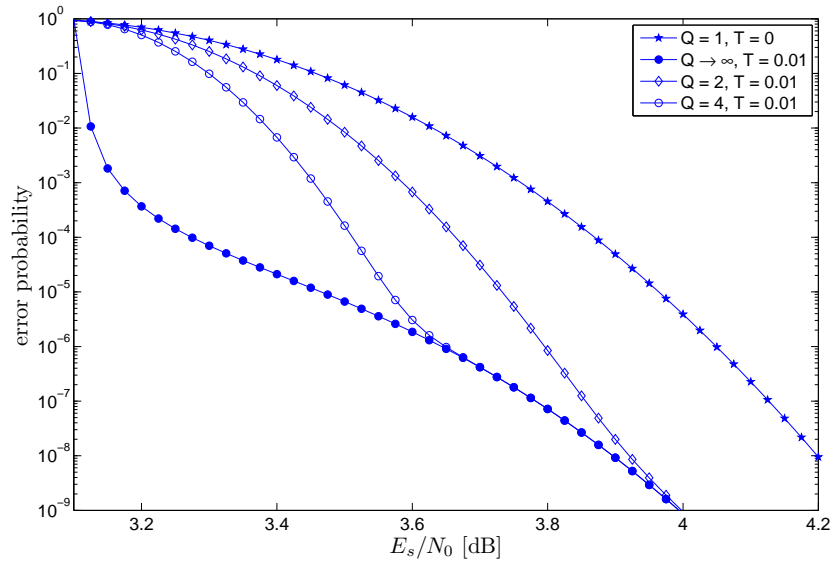
$E_0$  is as defined in (15), and  $q_X$  is an arbitrary probability distribution over  $\mathcal{X}$ .

*Proof:* The proof follows the same arguments as the proof of Corollary 1. ■

The following bound is provided for the case of binary linear block codes whose transmission takes place over an MBIOS channel (the generalization of the bound to non-binary linear block codes, as provided in [22], is direct):



(a) Lower bounds on the expected rates



(b) Upper bounds on the error probability

Fig. 4: Performance bounds of hybrid-ARQ schemes based on an expurgated, octal-alphabet and regular (8,16) LDPC code ensemble with a block length of  $n = 1008$  symbols. The transmission is assumed to take place over an AWGN channel with 8-PSK modulation. In plot (a), lower bounds on the expected rates for memoryless hybrid-ARQ schemes with and without deadlines (see (77), and (75), respectively) are shown for  $T = 0.01$  (and possible deadlines of  $Q = 2$  and 4 transmissions). In plot (b) upper bounds on the error probability are provided for the considered schemes.

**Corollary 5.** Consider an  $(n, k)$  binary linear block code  $\mathcal{C}$  whose transmission takes place over an MBIOS channel with a transition probability law  $p$ . Then the block error probability  $P_e$ , and the undetected error probability  $P_{ue}$ , under the generalized decoding rule in (7) satisfy

$$P_e \leq e^{-n\left(E(\rho, R, \mathcal{C}) - \frac{\rho T}{1+\rho}\right)}, \quad 0 \leq \rho \leq 1 \tag{37}$$

$$P_{ue} \leq e^{-n\left(E(\rho, R, \mathcal{C}) + \frac{\rho T}{1+\rho}\right)}, \quad 0 \leq \rho \leq 1 \tag{38}$$

where  $R$  is the code rate (in nats per channel use), and  $E(\rho, R, \mathcal{C})$  is defined in (21).

*Proof:* The proof follows from Proposition 7, and its derivation is similar to the way where Corollary 2 is derived from Proposition 6. ■

**Remark 15.** As in Corollary 2, the bounds of Corollary 5 resemble to the SFB, and they may therefore be considered as a generalization of the SFB for the case at hand.

**Remark 16.** For all rates below some (finite) rate thresholds, the bounds in Corollary 5 on the decoding error for linear block codes under the suboptimal LR rule in Definition 3, coincide with those under the optimal decoding rule in Definition 2. To see this, observe first that the upper bounds in (19) and (37) are identical. It is left to consider the upper bounds in (20) and (38) on the undetected error probability. Note first that  $E_0(\rho) - \rho R$  ( $E_0$  is defined in (22)) is a concave function of  $0 \leq \rho \leq 1$ , and it is optimized for rates below  $E'_0(1)$  at  $\rho = 1$  (see, e.g., [34, p. 135]). Moreover,  $\frac{\rho}{1+\rho}$  is a monotonic increasing function of  $0 \leq \rho \leq 1$ . This implies that if  $\frac{T}{4} < E'_0(1)$ , then at all rates below  $E'_0(1) - \frac{\ln(\alpha(\mathcal{C}))}{n} - \frac{T}{4}$ , the error exponents of the upper bounds in (20) and (38) are both maximized at  $\rho = 1$ , and they therefore coincide. A similar observation is provided in [21, p. 82] for the ensemble of fully random block codes. Specifically, it is observed in [21] that up to some rate threshold, the upper bounds under the suboptimal LR decoding rule for the ensemble of fully-random block codes coincide exponentially with those provided by Forney in [15].

**Example 7 (Error exponents of fully random binary linear block codes).** Fully random binary and linear  $(n, k)$  block codes are considered where, as mentioned in Example 3,  $\alpha(\mathcal{C}) = 1$  (see (23)). For the particular case of transmission over a BSC, the error exponents for the considered ensemble are studied in [2] and [3]. The lower bounds on the block error exponents and the undetected error exponents from [2] and [3] are compared in Fig. 5(a), and 5(b), respectively, to the bounds provided in Corollary 5. The bounds are derived for a BSC with a crossover probability of  $p = 0.07$  and a decoding parameter  $\tau = 0.03$  (see (8) where these are the same parameters studied in [2, Fig. 1]). The error exponent provided by Gallager for the case of ML decoding is also provided for comparison, in addition to the undetected error exponent under the optimal generalized decoding rule. Apart from low rates, where the bounds in [2] and [3] outperform those provided in Corollary 5, the latter bounds on the error exponents lie in between the two previously reported bounds from [2] and [3] (see Fig. 5). Moreover, in the rate region beyond the critical rate, where the bound in [2] outperform the bound in [3], the derived bounds perform in close proximity to the tightest known bound. The superiority of the undetected error exponent under the optimal decoding rule is clearly pronounced. As observed in Remark 16, it is evident that for low to moderate code rates, the bounds under optimal and suboptimal generalized decoding rules coincide. However, as the coding rates approach the channel capacity, the lower bounds on the undetected block error exponents under the suboptimal generalized-decoding, are considerably loosened in comparison to the lower bound under the optimal generalized decoding.

**Corollary 6.** Under the assumptions and notation in Corollary 3, the block error probability  $P_e$  and the undetected error probability  $P_{ue}$  under the suboptimal decoding rule in (7), satisfy

$$P_e \leq e^{\frac{n\rho T}{1+\rho}} \cdot D_s(\rho, \mathcal{C}), \quad 0 \leq \rho \leq 1 \quad (39)$$

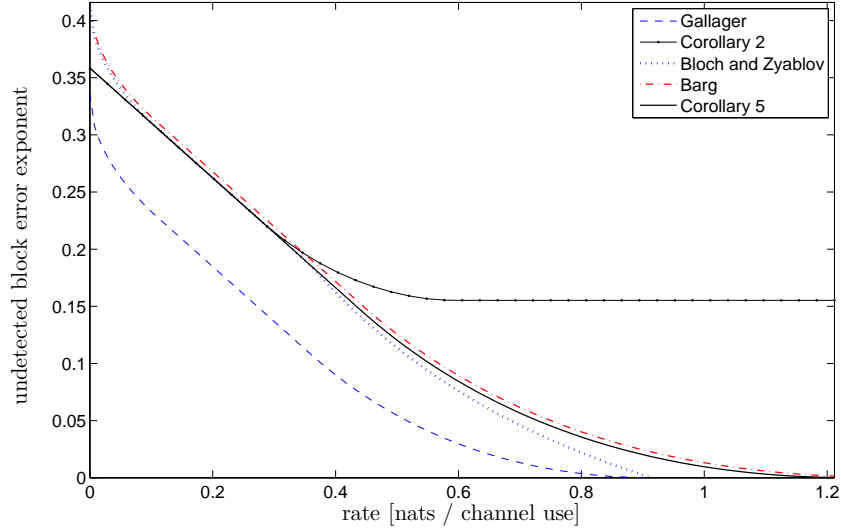
$$P_{ue} \leq e^{-\frac{n\rho T}{1+\rho}} \cdot D_s(\rho, \mathcal{C}), \quad 0 \leq \rho \leq 1 \quad (40)$$

where  $D_s(\rho, \mathcal{C})$  is defined in (29).

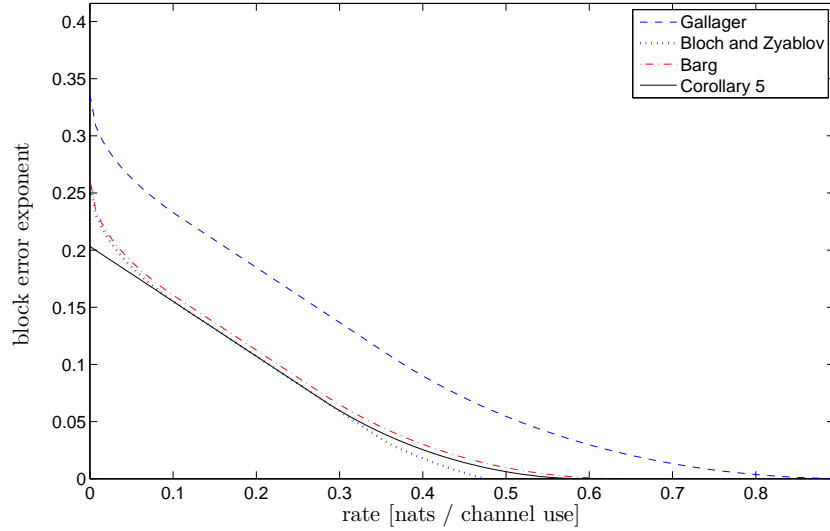
*Proof:* Setting  $s = \frac{\rho}{1+\rho}$ ,  $G_n^m(\mathbf{y}) = \prod_{i=1}^n g(y_i)$  where  $g$  is as defined in (24), the proof follows from Proposition 7 in the same way as the proof in [22, Theorem 3]. ■

Consider the particular case of binary linear block codes whose transmission takes place over the binary-input AWGN channel with BPSK modulation. The bound of Divsalar (see [9] and [26, Sec. 3.2.4]) provides a closed-form expression for an upper bound on the block error probability under ML decoding. The following proposition provides a similar bound under the LR decoding rule in Definition 3:

**Proposition 8.** Consider the transmission of a binary linear block code over the AWGN channel with BPSK



(a) Undetected error exponents



(b) Error exponents

**Fig. 5:** Lower bounds on the block error exponents of fully-random binary linear block codes whose transmission takes place over a BSC with a crossover probability of  $p = 0.07$ , under the suboptimal decoding rule in (8) with  $\tau = 0.03$ . The lower bounds on the undetected block error exponents in [2, Theorem 2], [3] (see also [2, Theorem 1]), and Corollary 5 (see (38)) are provided in plot (a), together with Gallager’s random-coding error exponent under ML decoding [17], and the lower bound on the undetected error exponent in Corollary 2 (see (20)) under the optimal generalized decoding rule. The lower bounds on the error exponents in [2, Theorem 2], [3], and Corollary 5 (see (37)) are provided in plot (b) (the lower bound of Gallager for the random-coding error exponent under ML decoding is also provided for comparison).

modulation, then the error and undetected error probabilities under the LR decoding in (7) satisfy

$$P_e \leq \sum_{d=d_{\min}}^n \min \left\{ \exp \left( -nE_e \left( \frac{d}{n}, \frac{E_s}{N_0} \right) \right), |C_d| Q \left( \sqrt{\frac{2E_s d}{N_0}} - \frac{nT}{2\sqrt{\frac{2dE_s}{N_0}}} \right) \right\} \quad (41)$$

$$P_{ue} \leq \sum_{d=d_{\min}}^n \min \left\{ \exp \left( -nE_{ue} \left( \frac{d}{n}, \frac{E_s}{N_0} \right) \right), |C_d| Q \left( \sqrt{\frac{2E_s d}{N_0}} + \frac{nT}{2\sqrt{\frac{2dE_s}{N_0}}} \right) \right\} \quad (42)$$

where  $d_{\min}$  is the minimum Hamming distance of the code,  $n$  is the block length of the code,  $|\mathcal{C}_i|$  is the number of codewords whose Hamming weight equals  $i$ ,  $T$  is the decoding parameter in (7),  $E_s$  is the energy per transmitted (coded) symbol,  $\frac{N_0}{2}$  is the two-sided power spectral density of the white Gaussian noise, and

$$\begin{aligned} E_e \left( \delta, \frac{E_s}{N_0} \right) &\triangleq E_D \left( \delta, \frac{E_s}{N_0} \right) - \frac{T\xi}{2}, \\ E_{ue} \left( \delta, \frac{E_s}{N_0} \right) &\triangleq E_D \left( \delta, \frac{E_s}{N_0} \right) + \frac{T\xi}{2} \\ E_D \left( \delta, \frac{E_s}{N_0} \right) &\triangleq -r_n(\delta) + \frac{1}{2} \ln \left( \beta + (1 - \beta)e^{2r_n(\delta)} \right) + \frac{\beta\delta}{1 - (1 - \beta)\delta} \frac{E_s}{N_0} \\ \beta &\triangleq \sqrt{\frac{E_s}{N_0} \frac{2(1 - \delta)}{\delta(1 - e^{-2r_n(\delta)})} + \left( \frac{1 - \delta}{\delta} \right)^2 \left( \left( 1 + \frac{E_s}{N_0} \right)^2 - 1 \right) - \frac{1 - \delta}{\delta} \left( 1 + \frac{E_s}{N_0} \right)} \\ r_n(\delta) &\triangleq \frac{\ln |\mathcal{C}_d|}{n}, \quad \delta \triangleq \frac{d}{n} \\ \xi &\triangleq \frac{\beta}{\beta + (1 - \beta)(1 - \delta)}. \end{aligned}$$

*Proof:* See Appendix I. ■

**Example 8 (Error performance of expurgated binary and regular LDPC code ensembles under suboptimal generalized decoding with erasures).** Consider an expurgation of the binary and regular LDPC code ensembles in Example 4 (with block lengths of 504 and 2004 bits). The upper bound in (40), on the undetected error probability under the generalized decoding rule with erasures in (7), is provided in Fig. 6 assuming that the transmission takes place over a binary-input AWGN channel. The upper bounds under the optimal generalized decoding rule are also provided for a comparison, in addition to the upper bound under the generalized decoding rule with  $T = 0$  (which coincides with the upper bound on the error probability under ML decoding). It is evident that the resulting bounds under the suboptimal generalized decoding rule are loosened in comparison to the bounds under the optimal generalized decoding rule. This result is expected from the previous example where the undetected error exponents are studied for fully-random linear block codes. In Fig 7, the upper bounds on the undetected error probability in Corollary 6 are compared with those provided in Proposition 8. The provided bounds are for the binary regular and expurgated LDPC code ensembles in Example 4 (with block lengths of 504 and 2004 bits), and for a similar ensemble with a block length of 10008 bits and  $D_n = 800$ . The parameter  $T$  in (7) is chosen, for this comparison, to be 0.0198, 0.0050, and  $9.992 \cdot 10^{-4}$ , respective to the considered block lengths. It is evident that the simple bound in (42) is loosened in comparison to the bound in (40), but only by a relatively small difference.

## V. UPPER BOUNDS UNDER FIXED-SIZE LIST DECODING

In this section, upper bounds on the block error probability are derived for the fixed-size list decoding (see Definition 4). As mentioned in Section II, the block error event in this case corresponds to the possibility that the decoded list does not include the transmitted codeword.

**Proposition 9.** Consider the transmission of a block code  $\mathcal{C}$  with  $M$  codewords of length  $n$ , and let  $p(\mathbf{y}|\mathbf{x})$  designate the transition probability of the channel where  $\mathbf{x} \in \mathcal{C}$  is the transmitted codeword and  $\mathbf{y} \in \mathcal{Y}^n$  is the received vector. Consider the case where a fixed-size list decoder is used where the size of the list is denoted by  $L$ . Then, the conditional block error probability  $P_{e|m}$ , given that the  $m$ -th message is transmitted satisfies

$$\begin{aligned} P_{e|m} &\leq \left( \sum_{\mathbf{y}} G_n^m(\mathbf{y}) p(\mathbf{y}|\mathbf{x}_m) \right)^{1-\rho} \\ &\quad \left( \frac{1}{L} \sum_{m' \neq m} \sum_{\mathbf{y}} p(\mathbf{y}|\mathbf{x}_{m'}) G_n^m(\mathbf{y})^{1-\frac{1}{\rho}} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{\rho}{\rho-1}} \right)^{\rho}. \end{aligned} \quad (43)$$



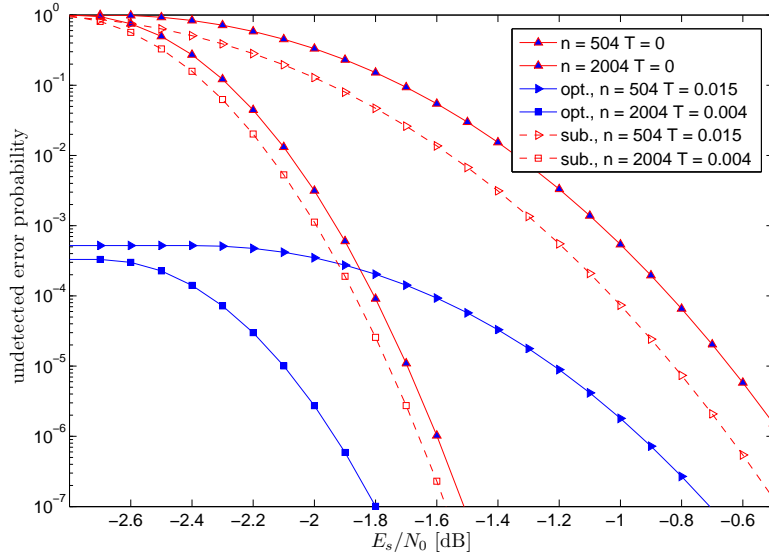


Fig. 6: Upper bounds on the undetected error probabilities of some expurgated ensembles of binary and regular (6,12) LDPC codes under the optimal and sub-optimal generalized decoding rules in (5) and (7), respectively. The upper bound in Corollary 6 is evaluated assuming that the transmission takes place over a binary-input AWGN channel. The upper bounds in Corollary 3, studied in Examples 4 and 5, are also provided for comparison.

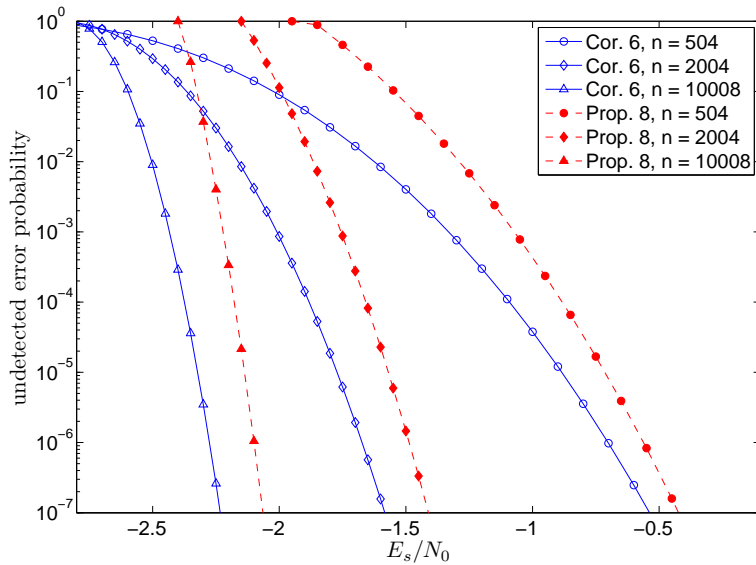


Fig. 7: A comparison between the upper bounds in (40) and (42), on the undetected error probability under the LR generalized decoding rule in (7). The comparison is provided for binary expurgated and regular (6,12) LDPC code ensembles of Gallager with block lengths of 504, 2004 and 10008 bits whose transmissions take place over binary-input AWGN channels with BPSK modulation.

where  $0 \leq s \leq \rho \leq 1$  are real-valued parameters, and  $G_n^m$  is an arbitrary non-negative function over  $\mathcal{Y}^n$  which possibly depends on the codeword  $\mathbf{x}_m$ , for  $1 \leq m \leq M$ .

*Proof:* See Appendix J. ■

The following corollary is a particularization of Proposition 9 for the ensemble of fully-random block codes, with fixed block length and rate, whose transmission takes place over a memoryless channel:

**Corollary 7.** Consider the transmission of a block code  $\mathcal{C}$  over a memoryless communication channel. Then, under

the notation in Proposition 9, there exists a block code whose block error probability  $P_e$  under fixed-size list decoding satisfies

$$P_e \leq e^{-nE_r(R - \frac{1}{n} \ln L)} \quad (44)$$

where  $R \triangleq \frac{\ln M}{n}$  is the code rate (in nats per channel use),

$$E_r(R) \triangleq \max_{0 \leq \rho \leq 1, q_X} \left( E_0(\rho, q_X) - \rho R \right) \quad (45)$$

$$E_0(\rho, q_X) \triangleq -\ln \left( \sum_{y \in \mathcal{Y}} \left( \sum_{x \in \mathcal{X}} q_X(x) p(y|x)^{\frac{1}{1+\rho}} \right)^{1+\rho} \right)$$

and  $q_X$  is a probability distribution over the input alphabet  $\mathcal{X}$ .

*Proof:* Fix a probability distribution  $q_X$  over  $\mathcal{X}$ , and consider the ensemble of random block codes where each codeword is chosen independently according to  $q_{\mathbf{X}}(\mathbf{x}) = \prod_{i=1}^n q_X(x_i)$ . First, we apply the bound in (43) for a specific realization of a codebook, with  $s = \frac{\rho}{1+\rho}$  and

$$G_n^m(\mathbf{y}) \triangleq \left( \sum_{\mathbf{x}} q_{\mathbf{X}}(\mathbf{x}) \left( \frac{p(\mathbf{y}|\mathbf{x})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^{\rho}.$$

The proof follows by a random coding argument, and by choosing the optimal probability distribution  $q_X$ . ■

**Remark 17 (On comparison of the error exponent in Corollary 7 with previously known results).** The upper bound in Corollary 7 is compared to three previously known results:

- 1) **The sphere-packing bound:** The sphere-packing lower bound in [28, eq. 1.6] provides an exponential lower bound on the error probability for fixed-size list-decoding of block codes. The bound in Corollary 7 and the sphere-packing bound exponentially coincide for all rates above the critical rate (where the maximization of the random coding error exponent is achieved for  $0 \leq \rho \leq 1$ ).
- 2) **Asymptotic upper bound:** Consider the case where the size of the decoded list grows exponentially with the blocklength, and denote the exponential growth rate of the decoded list by  $l$  (i.e.,  $L = e^{nl}$  for some  $l > 0$ ). The following asymptotic upper bound is provided in [8, p. 196, ex. 27] for the case at hand:

$$\limsup_{n \rightarrow \infty} \frac{1}{n} \ln P_e \leq -E_r(R - l). \quad (46)$$

It is easily verified that the bound in Corollary 7 asymptotically coincides with the bound in (46).

- 3) **A variation on the Gallager bound:** The following exponential upper bound on the error probability is provided in [18, p. 538, ex. 5.20] for given block length and list size (the same assumptions and notation as in Corollary 7 are considered):

$$P_e \leq e^{-nE_r(R, L)}$$

where

$$E_r(R, L) \triangleq \max_{0 \leq \rho \leq L, q_X} \left( E_0(\rho, q_X) - \rho R \right). \quad (47)$$

The error exponents in (44) and (47) differ in the following aspects:

- a) For a fixed list-size  $L$ , the error exponent in (44) depends on the block length  $n$  while the error exponent in (47) does not.
- b) The maximization of  $\rho$  in (44) is carried over the interval  $[0, 1]$  while in (47) it is  $[0, L]$ .
- c) The bound in (44) includes an explicit rate reduction term, which depends on the list size.
- d) The derivation of the bound in (44) is based on a particularization of the DS2 bound in Proposition 9 for fully-random block codes. On the other hand, the derivation of the bound in (47) is based on a modification of the random coding bound [17] for the case at hand.

The two bounds in (44) and (47) are compared in Figure 8. Transmission of fully-random block codes over a BSC with a crossover of  $p = 0.11$  are considered, where equiprobable  $q_X(x) = \frac{1}{2}$ ,  $x \in \mathcal{X}$ , is assumed. The error exponent  $E_r(R, L)$  in (47) is plotted for a list size of  $L = 16$  codewords. In addition, the exponent

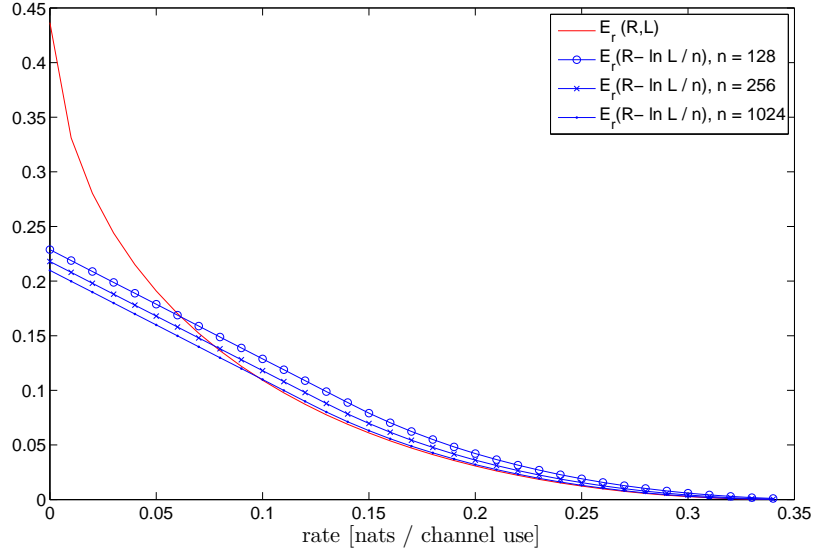


Fig. 8: A comparison between the upper bounds in Corollary 7 and [18, p. 538, ex. 5.20]. Transmission of a fully-random binary block codes (with independent equiprobable selection of coded bits) over a BSC with a cross over probability of  $p = 0.11$  is assumed. The exponent term  $E_r(R, L)$  in (47) is plotted for a list size of  $L = 16$  codewords. The exponent  $E_r(R - \frac{1}{n} \ln L)$  in (44) is plotted for the same list-size and blocklengths of 128, 256 and 1024 bits.

$E_r(R - \frac{1}{n} \ln L)$  is provided for the same list size and block lengths of 128, 256 and 1024 bits. It is observed that for low rates the bound in (47) outperforms the bound in (44). For moderate rates, the bound in (44) outperforms the bound in (47). The gap between the plotted exponents is negligible as the block length increases (even for a moderate block length of 1024 bits).

The following bound is provided for the case of binary linear block codes whose transmission takes place over an MBIOS channel:

**Corollary 8.** Consider an  $(n, k)$  binary linear block code  $\mathcal{C}$  whose transmission takes place over an MBIOS channel. Then, the block error probability  $P_e$  under fixed-size list-decoding, satisfies

$$P_e \leq e^{-nE_r\left(R + \frac{1}{n} \ln\left(\frac{\alpha(\mathcal{C})}{L}\right)\right)} \quad (48)$$

where

$$E_r(R) \triangleq \max_{0 \leq \rho \leq 1} (E_0(\rho) - \rho R)$$

and  $R$  is the code rate (in nats per channel use),  $L$  is the list size, and  $E_0(\rho)$  and  $\alpha(\mathcal{C})$  are defined in (22) and (23), respectively.

*Proof:* According to Proposition 4, it is necessary to analyze only the conditional error event assuming that the all-zero codeword is transmitted. Setting  $G_n^0(\mathbf{y}) = \prod_{i=1}^n g(y_i)$  in (43), it follows that

$$P_e \leq \left( \sum_{y \in \mathcal{Y}} g(y) p(y|0) \right)^{n(1-\rho)} \left( \frac{1}{L} \sum_{i=1}^n |\mathcal{C}_i| \left( \sum_{y \in \mathcal{Y}} g(y)^{1-\frac{1}{\rho}} p(y|0) \right)^{n-i} \left( \sum_{y \in \mathcal{Y}} g(y)^{1-\frac{1}{\rho}} p(y|1)^\lambda p(y|0)^{1-\lambda} \right)^i \right)^\rho \quad (49)$$

where  $|\mathcal{C}_i|$  denotes the number of codewords whose Hamming distance is  $i$ ,  $1 \leq i \leq n$ . The proof follows from (49) by setting  $\lambda = \frac{1}{1+\rho}$  where  $g$  is as defined in (24) (see similar derivation in [26, Section 4.4.1]). ■

**Remark 18.** For the particular case of fully-random linear block codes, the bound in (48) coincides with the bound in Corollary 7 for fully-random block codes.

**Remark 19.** The bound in Corollary 8 resembles to the SFB [29], and therefore may be considered as a generalization of the SFB for the case at hand.

**Remark 20.** The bound in (49) can be generalized to non-binary linear block codes using a similar derivation as in [22]. Note, however, that in [22], non-binary codes are studied under ML decoding and not list-decoding. Nevertheless, the similarity of the bound in (43) to the upper bounds derived in [22] allows to use the same arguments for the case at hand (see Appendix J).

**Corollary 9.** Under the assumptions and notation in Corollary 3, the block error probability  $P_e$  under fixed-size list-decoding where  $L$  denotes the size of the list, satisfies

$$P_e \leq A(\rho)^{n(1-\rho)} \left( \frac{1}{L} \sum_{1 \leq l \leq n} P(l) \binom{n}{l} B(\rho)^{n-l} C(\rho)^l \right)^\rho \quad (50)$$

where  $A(\rho)$ ,  $B(\rho)$ , and  $C(\rho)$  are defined in (30)–(32).

*Proof:* Setting  $s = \frac{\rho}{1+\rho}$  and  $G_n^m(\mathbf{y}) = \prod_{i=1}^n g(y_i)$  where  $g$  is defined in (24), the proof follows from Proposition 9 in the same way as the proof in [22, Theorem 3]. ■

**Remark 21.** In the derivation of the bound in (48), a sum is upper bounded by a product of the maximal summand with the number of summands. This operation is avoided in the derivation of the bound in (50). Hence, the bound in Corollary 9 is tighter than the one in Corollary 8.

**Remark 22.** For the particular case of binary linear block codes, the symmetry condition in (26) is not mandatory and the bound in Corollary 9 follows by replacing the term  $P(l) \binom{n}{l}$  with the distance spectrum of the considered code (ensemble).

**Example 9 (Error performance of an expurgated ensemble of regular LDPC codes under fixed-size list decoding).** Consider the expurgation of Gallager’s ensemble of binary and regular (6,12) LDPC codes with a block length of 2004 bits (see Example 4). Upper bounds on the block error probability under fixed-size list-decoding are shown in Fig. 9 assuming that the transmission takes place over a binary-input AWGN channel. The upper bound in Corollary 9 is evaluated for list sizes of  $L = 1, 16,$  and  $128$  codewords. Note that the upper bound for  $L = 1$  corresponds to ML decoding. The bounds on the error probability show some marginal improvement by increasing the considered list size from  $L = 1$  to  $128$ . Similar conclusions can be observed for non-binary ensembles.

## VI. SUMMARY AND CONCLUSIONS

This paper considers performance bounds for several generalized decoding rules over memoryless symmetric channels. Three types of generalized decoding rules are considered:

- 1) The optimal generalized decoding rule in [15] with erasures and variable list sizes.
- 2) The suboptimal likelihood-ratio (LR) decoding rule with erasures (see [2] and [15]).
- 3) A fixed-size list decoding rule (see [14] and [36]) where the decoder outputs a list which includes the  $L$  most probable codewords (where the value of  $L$  is set a-priori).

The independence of the error performance on the transmitted codeword is proved in Propositions 2-4 for the considered decoding rules.

Upper bounds on the decoding error probability are provided. These bounds are suitable for the analysis of structured and random codes (or code ensembles) over memoryless symmetric channels. Both binary and non-binary code ensembles are studied in this paper under generalized decoding rules. When binary codes are considered, the bounds are based on the distance spectra of the codes, and when non-binary ensembles are studied, the complete composition spectra are required under the symmetry assumption in (26). For the case of LR decoding of binary linear block codes, a derivation of a closed-form expression is provided via a similar derivation to [9] which applies to ML decoding.

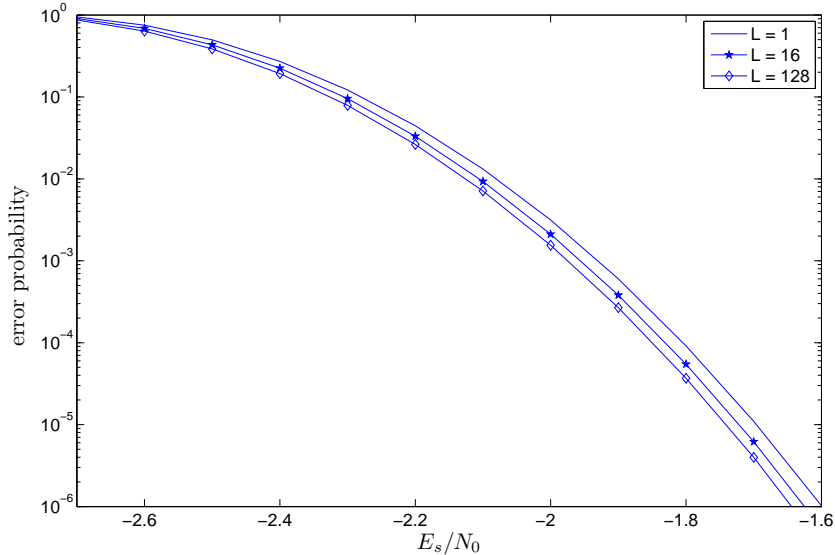


Fig. 9: Upper bounds on the error probability for an expurgation of Gallager’s ensemble of binary and regular (6,12) LDPC codes with a block length of 2004 bits (see Example 4). A list decoder is assumed where the size of the list is set to  $L$ . The upper bound in Corollary 9 is provided for some values of  $L$ , where the transmission takes place over a binary-input AWGN channel.

Several particular cases of the provided bounds are studied. The random coding error exponents in [15] are reproduced. In addition, error exponents under the suboptimal LR decoding rule with erasures are also derived. These error exponents are derived by applying the new bounds to fully random block codes. Next, a derivation of the error exponents of fully random linear block codes under optimal and suboptimal (LR) generalized decoding is provided. The resulting error exponents under the suboptimal LR decoding rule are compared with a recent improvement in [2], where the ensemble of binary fully random linear block codes over binary symmetric channels (BSC) is studied. This comparison shows good match with the provided error exponents with the results in [2]. In addition, it is shown that the error exponents for the fully random linear block codes under the suboptimal LR decoding rule, coincide for low rates with the corresponding error exponents under the optimal decoding rule. This is similar to an observation in [21], where the ensemble of fully random block codes is considered. A lower bound on the error exponent under fixed-size list-decoding is also studied as an application. This bound is compared to the sphere-packing lower bound on the error probability [28], and two additional upper bounds on the error probability, provided in [8] and [18].

Applications of the bounds for the performance analysis of structured code ensembles are further exemplified for some expurgated ensembles of (binary and non-binary) regular low-density parity-check (LDPC) codes. The error performance under some generalized decoding rules for these LDPC code ensembles is studied assuming that the transmission takes place over memoryless symmetric channels. The application of the provided bounds for the study of hybrid automatic-repeat request (ARQ) schemes is also demonstrated. The possibility of further investigating and optimizing the trade-offs between undetected error and erasures is suggested for further study in the context of linear block codes, based on the derived bounds.

#### APPENDIX A PROOF OF PROPOSITION 2

The following proof holds for memoryless symmetric channels with discrete-output alphabets, and the generalization to continuous-output alphabets is direct. We state first the following technical lemma:

**Lemma 1.** let  $x_1, x_2, x_3$  be arbitrary symbols in  $\mathcal{X}$ , and let  $p$  be a transition probability law of a memoryless symmetric channel. Then,

$$p\left(\mathcal{T}(\mathcal{T}(y, x_1), x_2) | x_3\right) = p(\mathcal{T}(y, x_1 + x_2) | x_3)$$

where  $\mathcal{T}$  is a mapping which satisfies the properties in Definition 1.

*Proof:* the reader is referred to [22, Appendix A]. ■

Assuming that all the codewords are sent with equal probability, the decision regions in (5) satisfy

$$\begin{aligned} \Lambda_m &\stackrel{(a)}{=} \left\{ \mathbf{y} : \frac{p(\mathbf{y}|\mathbf{x}_m)}{\sum_{m' \neq m} p(\mathbf{y}|\mathbf{x}_{m'})} \geq e^{nT} \right\} \\ &\stackrel{(b)}{=} \left\{ \mathbf{y} : \frac{\prod_{i=1}^n p(y_i|x_{m,i})}{\sum_{m' \neq m} \prod_{i=1}^n p(y_i|x_{m',i})} \geq e^{nT} \right\} \\ &\stackrel{(c)}{=} \left\{ \mathbf{y} : \frac{\prod_{i=1}^n p(\mathcal{T}(y_i, -x_{m,i})|0)}{\sum_{m' \neq m} \prod_{i=1}^n p(\mathcal{T}(y_i, -x_{m',i})|0)} \geq e^{nT} \right\} \end{aligned} \quad (51)$$

where (a) follows from (5) and the equal a-priori message probability assumption, (b) holds since the channel is memoryless, and (c) follows from the symmetry of the channel (see (1)). Let  $\mathbf{z} = (z_1, \dots, z_n)$  be defined as

$$z_i \triangleq \mathcal{T}(y_i, -x_{m,i}), \quad 1 \leq i \leq n \quad (52)$$

where  $m$  is the index of the transmitted codeword. From Lemma 1, it follows that  $\mathbf{y} \in \Lambda_m$  if and only if  $\mathbf{z} \in \tilde{\Lambda}_m$  where

$$\tilde{\Lambda}_m \triangleq \left\{ \mathbf{z} \in \mathcal{Y}^n : \frac{\prod_{i=1}^n p(z_i|0)}{\sum_{m' \neq m} \prod_{i=1}^n p(\mathcal{T}(z_i, x_{m,i} - x_{m',i})|0)} \geq e^{nT} \right\}, \quad 1 \leq m \leq q^k.$$

Using the linearity of the code, it follows that

$$\tilde{\Lambda}_m = \left\{ \mathbf{z} \in \mathcal{Y}^n : \frac{\prod_{i=1}^n p(z_i|0)}{\sum_{l \neq 0} \prod_{i=1}^n p(\mathcal{T}(z_i, x_{l,i})|0)} \geq e^{nT} \right\}.$$

Since the set  $\tilde{\Lambda}_m$  is independent of the index  $m$ , then

$$\tilde{\Lambda}_m = \tilde{\Lambda}_1 \text{ for all } 1 \leq m \leq q^k. \quad (53)$$

As a result, the conditional block error probability of the  $m$ -th message in (2) satisfies

$$\begin{aligned} P_{e|m} &= \sum_{\mathbf{z} \in \tilde{\Lambda}_m^c} p(\mathbf{z}|\mathbf{0}) \\ &\stackrel{(a)}{=} \sum_{\mathbf{z} \in \tilde{\Lambda}_1^c} p(\mathbf{z}|\mathbf{0}) \end{aligned}$$

where (a) follows from (53). This concludes the proof of the message independence property for the block error event.

We continue in proving the message independence property for the undetected error event (or the expected number of incorrect codewords when list decoding is considered). Assuming a memoryless symmetric channel, it follows from (1) and (4) that

$$\begin{aligned} P_{ue|m} &= \sum_{m' \neq m} \sum_{\mathbf{y} \in \Lambda_{m'}} p(\mathbf{y}|\mathbf{x}_m) \\ &= \sum_{m' \neq m} \sum_{\mathbf{y} \in \Lambda_{m'}} \prod_{i=1}^n p(\mathcal{T}(y_i, -x_{m,i})|0) \end{aligned} \quad (54)$$

where from (51)

$$\Lambda_{m'} = \left\{ \mathbf{y} : \frac{\prod_{i=1}^n p(\mathcal{T}(y_i, -x_{m',i})|0)}{\sum_{m'' \neq m'} \prod_{i=1}^n p(\mathcal{T}(y_i, -x_{m'',i})|0)} \geq e^{nT} \right\}.$$

Let  $\mathbf{z}$  be a vector defined as in (52), then from Lemma 1

$$p(\mathcal{T}(y_i, -x_{m',i})|0) = p(\mathcal{T}(z_i, x_{m,i} - x_{m',i})|0), \quad i = 1, \dots, n.$$

Hence, given that  $\mathbf{x}_m$  is the transmitted codeword, then  $\mathbf{y} \in \Lambda_{m'}$  for some  $m' \neq m$  if and only if  $\mathbf{z} \in \Gamma_{m,m'}$  where

$$\Gamma_{m,m'} \triangleq \left\{ \mathbf{z} \in \mathcal{Y}^n : \frac{\prod_{i=1}^n p(\mathcal{T}(z_i, x_{m,i} - x_{m',i})|0)}{\sum_{m'' \neq m'} \prod_{i=1}^n p(\mathcal{T}(z_i, x_{m,i} - x_{m'',i})|0)} \geq e^{nT} \right\}. \quad (55)$$

From (52), the conditional undetected error probability in (54) is rewritten in the form

$$P_{\text{ue}|m} = \sum_{m' \neq m} \sum_{\mathbf{z} \in \Gamma_{m,m'}} p(\mathbf{z}|\mathbf{0}). \quad (56)$$

Using the linearity of the code, then  $x_{m,i} - x_{m'',i} = (x_{m,i} - x_{m',i}) + (x_{m',i} - x_{m'',i}) = x_{l_1,i} + x_{l_2,i}$  for some indices  $l_1$  and  $l_2$  which correspond to non-zero codewords. Let  $\mathbf{x} \triangleq \mathbf{x}_{l_1}$  and  $\tilde{\mathbf{x}} = \mathbf{x}_{l_2}$ , then the conditional undetected error probability in (56) is expressed equivalently in the form

$$P_{\text{ue}|m} = \sum_{\substack{\mathbf{x} \in \mathcal{C} \\ \mathbf{x} \neq \mathbf{0}}} \sum_{\mathbf{z} \in \Gamma(\mathbf{x})} p(\mathbf{z}|\mathbf{0})$$

where, based on (55),

$$\Gamma(\mathbf{x}) \triangleq \left\{ \mathbf{z} \in \mathcal{Y}^n : \frac{\prod_{i=1}^n p(\mathcal{T}(z_i, x_i)|0)}{\sum_{\substack{\tilde{\mathbf{x}} \in \mathcal{C} \\ \tilde{\mathbf{x}} \neq \mathbf{0}}} \prod_{i=1}^n p(\mathcal{T}(z_i, x_i + \tilde{x}_i)|0)} \geq e^{nT} \right\}.$$

This proves the independence property for the undetected error event, and it concludes the proof of Proposition 2.

## APPENDIX B PROOF OF PROPOSITION 3

Similarly to Appendix A, also the following proof considers memoryless symmetric channels with discrete-output alphabets, where the generalization to continuous output alphabets is direct. Let  $p$  be the transition probability function of the considered channel,  $\mathcal{C}$  be an  $(n, k)$  linear block code over an alphabet whose cardinality is  $q$ , and  $\mathcal{T}$  be a mapping as specified in Definition 1. It is assumed that all the codewords of  $\mathcal{C}$  are sent with equal probability. For an arbitrary set  $\Lambda \subseteq \mathcal{Y}^n$  and a codeword  $\mathbf{x}_m \in \mathcal{C}$ , let

$$\mathcal{Z}_m(\Lambda) \triangleq \left\{ \mathbf{z} \in \mathcal{Y}^n : (\mathcal{T}(z_1, x_{m,1}), \mathcal{T}(z_2, x_{m,2}), \dots, \mathcal{T}(z_n, x_{m,n})) \in \Lambda \right\}. \quad (57)$$

In addition, we use the notation  $\Lambda^{\text{LR}}(\mathbf{x}_m)$  for the decision region  $\Lambda_m^{\text{LR}}$  in (7) of the codeword  $\mathbf{x}_m$ . Note that for the concerned decoding rule with  $T > 0$ , the decision regions are disjoint. The following technical lemma is introduced:

**Lemma 2.** Let  $\mathcal{Z}_m$  be the mapping defined in (57), and  $\Lambda_m^{\text{LR}}$  be the decision region in (7). Then,

$$\mathcal{Z}_m(\Lambda_{m'}^{\text{LR}}) = \Lambda^{\text{LR}}(\mathbf{x}_{m'} - \mathbf{x}_m), \quad \forall m, m' \in \{1, \dots, q^k\}. \quad (58)$$

*Proof:* Let us choose  $\mathbf{z} \in \mathcal{Z}_m(\Lambda_{m'}^{\text{LR}})$ , and let  $\mathbf{y} = (y_1, \dots, y_n)$  be defined via the equality

$$y_i = \mathcal{T}(z_i, x_{m,i}), \quad i = 1, \dots, n. \quad (59)$$

From (7) and (57)

$$\frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_{m'_2})} \geq e^{nT}$$

where  $\mathbf{x}_{m'}$  and  $\mathbf{x}_{m'_2}$  are the most probable codewords, in a descending order, for  $\mathbf{y}$  as a received vector. Using the symmetry of the channel, it follows from (1) that

$$p(\mathbf{y}|\mathbf{x}_{m'}) = p(\mathbf{z}|\mathbf{x}_{m'} - \mathbf{x}_m).$$

As a result,  $\mathbf{x}_{m'} - \mathbf{x}_m$  is the most probable codeword if  $\mathbf{z}$  is the received vector (otherwise, if there exists a codeword  $\mathbf{x} \neq \mathbf{x}_{m'} - \mathbf{x}_m$  which is more probable, then there exists a more probable codeword for  $\mathbf{y}$  which is different from  $\mathbf{x}_{m'}$ ). The same argument shows that  $\mathbf{x}_{m'_2} - \mathbf{x}_m$  is the second most probable codeword for  $\mathbf{z}$ , and

$$\frac{p(\mathbf{z}|\mathbf{x}_{m'} - \mathbf{x}_m)}{p(\mathbf{z}|\mathbf{x}_{m'_2} - \mathbf{x}_m)} \geq e^{nT}.$$

This verifies that  $\mathbf{z} \in \Lambda^{\text{LR}}(\mathbf{x}_{m'} - \mathbf{x}_m)$  which shows that  $\mathcal{Z}_m(\Lambda_m^{\text{LR}}) \subseteq \Lambda^{\text{LR}}(\mathbf{x}_{m'} - \mathbf{x}_m)$ . To show the opposite inclusion, which then yields that these two sets are equal, let  $\mathbf{z} \in \Lambda^{\text{LR}}(\mathbf{x}_{m'} - \mathbf{x}_m)$ . This implies that the codeword  $\mathbf{x}_{m'} - \mathbf{x}_m$  is the most probable codeword if  $\mathbf{z}$  is the received vector, and

$$\frac{p(\mathbf{z}|\mathbf{x}_{m'} - \mathbf{x}_m)}{p(\mathbf{z}|\mathbf{x}_{m_2'')}} \geq e^{nT}$$

where  $\mathbf{x}_{m_2''}$  is the second most probable codeword for  $\mathbf{z}$ . Again, using the symmetry of the channel, for a vector  $\mathbf{y}$  as in (59), it follows that  $\mathbf{x}_{m'}$  is the most probable codeword for  $\mathbf{y}$ ,  $\mathbf{x}_{m_2''} + \mathbf{x}_m$  is the second most probable codeword for  $\mathbf{y}$ , and

$$\frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_{m_2''} + \mathbf{x}_m)} \geq e^{nT}.$$

As a result,  $\mathbf{z} \in \mathcal{Z}_m(\Lambda_m^{\text{LR}})$ , which yields that  $\Lambda^{\text{LR}}(\mathbf{x}_{m'} - \mathbf{x}_m) \subseteq \mathcal{Z}_m(\Lambda_m^{\text{LR}})$ . This concludes the proof of (58). ■

From (59), the conditional block error probability satisfies

$$\begin{aligned} P_{e|m} &= \sum_{\mathbf{y} \notin \Lambda_m^{\text{LR}}} p(\mathbf{y}|\mathbf{x}_m) \\ &\stackrel{(a)}{=} \sum_{\mathbf{z} \notin \mathcal{Z}_m(\Lambda_m^{\text{LR}})} p(\mathbf{z}|\mathbf{0}) \\ &\stackrel{(b)}{=} \sum_{\mathbf{z} \notin \Lambda^{\text{LR}}(\mathbf{0})} p(\mathbf{z}|\mathbf{0}) \end{aligned}$$

where (a) follows from (1) and (59), and (b) follows from (58). This proves the message independence property for the conditional block error probability. Using the same arguments, the message independence property is established for the conditional undetected error probability:

$$\begin{aligned} P_{ue|m} &= \sum_{m' \neq m} \sum_{\mathbf{y} \in \Lambda_{m'}^{\text{LR}}} p(\mathbf{y}|\mathbf{x}_m) \\ &= \sum_{m' \neq m} \sum_{\mathbf{z} \in \mathcal{Z}_m(\Lambda_{m'}^{\text{LR}})} p(\mathbf{z}|\mathbf{0}) \\ &= \sum_{m' \neq m} \sum_{\mathbf{z} \in \Lambda^{\text{LR}}(\mathbf{x}_{m'} - \mathbf{x}_m)} p(\mathbf{z}|\mathbf{0}) \\ &= \sum_{\substack{\mathbf{x} \in \mathcal{C} \\ \mathbf{x} \neq \mathbf{0}}} \sum_{\mathbf{z} \in \Lambda^{\text{LR}}(\mathbf{x})} p(\mathbf{z}|\mathbf{0}) \end{aligned}$$

where the second equality follows from (59) and since the mapping  $\mathcal{T}$  is bijective, the third equality follows from (58), and the last equality follows from the linearity of the code.

#### APPENDIX C PROOF OF PROPOSITION 4

Considering ties as error events<sup>2</sup>, the conditional block error probability for a list of size  $L$  satisfies

$$P_{e|m} = \sum_{\mathbf{y} \in \Lambda_m^L} p(\mathbf{y}|\mathbf{x}_m) \tag{60}$$

where

$$\Lambda_m^L \triangleq \{\mathbf{y} \in \mathcal{Y}^n : \exists \{m_i\}_{i=1}^L \text{ s.t. } m_i \neq m, p(\mathbf{y}|\mathbf{x}_{m_i}) \geq p(\mathbf{y}|\mathbf{x}_m) \forall 1 \leq i \leq L\} \tag{61}$$

is the complementary of the decision region of  $\mathbf{x}_m \in \mathcal{C}$  under list decoding of fixed-size  $L$  (here  $\{m_i\}_{i=1}^L$  is a sequence of distinct integers), i.e., if  $\mathbf{y} \in \Lambda_m^L$  then the codeword  $\mathbf{x}_m$  is not included in the list for a received vector

<sup>2</sup>Such a pessimistic assumption is reasonable, see also a similar assumption in [34, p. 59].



$\mathbf{y}$ . Using the change of variables in (59), it follows from (60) that for linear block codes whose transmission takes place over memoryless symmetric channels

$$P_{e|m} = \sum_{\mathbf{z} \in \mathcal{Z}_m(\Lambda_m^L)} p(\mathbf{z}|\mathbf{0})$$

where  $\mathcal{Z}_m(\Lambda_m^L)$  is as defined in (57). The following lemma concludes the proof of Proposition 4:

**Lemma 3.** Let  $\mathcal{Z}_m$  be a mapping defined in (57), and  $\Lambda_m^L$  be the decoding region of  $\mathbf{x}_m \in \mathcal{C}$  under list decoding with a fixed size  $L$ . Then,

$$\mathcal{Z}_m(\Lambda_m^L) = \Lambda_1^L$$

for all  $1 \leq m \leq q^k$ , where  $\Lambda_1^L$  is the complementary of the decision region of the all-zero codeword  $\mathbf{x}_1 = \mathbf{0}$  under list decoding of size  $L$ .

*Proof:* Let us choose  $\mathbf{z} \in \mathcal{Z}(\Lambda_m^L)$ . From (57), there exists  $\mathbf{y} \in \Lambda_m^L$  where

$$y_i = \mathcal{T}(z_i, x_{m,i}), \quad i = 1, \dots, n \quad (62)$$

and  $\mathcal{T}$  is as specified in Definition 1. From (61), there exists a list of  $L$  distinct codewords,  $\{\mathbf{x}_{m_i}\}_{i=1}^L$ , for which

$$p(\mathbf{y}|\mathbf{x}_{m_i}) > p(\mathbf{y}|\mathbf{x}_m), \quad i = 1, \dots, L. \quad (63)$$

Using the symmetry of the channel, it follows that

$$p(\mathbf{z}|\mathbf{x}_{m_i} - \mathbf{x}_m) \geq p(\mathbf{z}|\mathbf{0}). \quad (64)$$

This assures that  $\mathbf{z} \in \Lambda_1^L$ , which shows that  $\mathcal{Z}_m(\Lambda_m^L) \subseteq \Lambda_1^L$ .

Next, in order to show the opposite inclusion, let  $\mathbf{z} \in \Lambda_1^L$ . Then, there exists a list of  $L$  non-zero codewords  $\{\mathbf{x}_{m_i}\}_{i=1}^L$ ,  $m_i \neq 1$ , satisfying

$$p(\mathbf{z}|\mathbf{x}_{m_i}) \geq p(\mathbf{z}|\mathbf{0})$$

and therefore from the symmetry of the mapping  $\mathcal{T}$  and the equality in (62), we get

$$p(\mathbf{y}|\mathbf{x}_{m_i} + \mathbf{x}_m) \geq p(\mathbf{y}|\mathbf{x}_m)$$

It assures that  $\mathbf{z} \in \mathcal{Z}_m(\Lambda_m^L)$  which implies that  $\Lambda_1^L \subseteq \mathcal{Z}_m(\Lambda_m^L)$ . This two inclusions complete the proof of the lemma. ■

#### APPENDIX D PROOF OF PROPOSITION 5

Let  $\Lambda_m$  be the generalized decision region as defined in (5). For  $\mathbf{y} \notin \Lambda_m$ , it follows that

$$1 = e^{nT} e^{-nT} \leq e^{nT} \left( \sum_{m' \neq m} \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right). \quad (65)$$

Let  $s$  and  $\rho$  satisfy  $0 \leq s \leq \rho \leq 1$ , and recall the following inequality (see [34, p.197]):

$$\sum_i a_i \leq \left( \sum_i a_i^\lambda \right)^{\frac{1}{\lambda}} \quad (66)$$

which holds if  $a_i \geq 0$  and  $0 < \lambda \leq 1$ . Setting

$$a_i = \frac{p(\mathbf{y}|\mathbf{x}_i)}{p(\mathbf{y}|\mathbf{x}_m)}, \quad \lambda = \frac{s}{\rho}$$

it follows from (2), (65) and (66) that the conditional error probability of the  $m$ -th message satisfies

$$\begin{aligned} P_{e|m} &\leq e^{nTs} \sum_{\mathbf{y} \in \Lambda_m^c} p(\mathbf{y}|\mathbf{x}_m) \left( \sum_{m' \neq m} \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^s \\ &\leq e^{nTs} \sum_{\mathbf{y} \in \Lambda_m^c} p(\mathbf{y}|\mathbf{x}_m) \left( \sum_{m' \neq m} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^\rho. \end{aligned} \quad (67)$$

Let  $\psi_n^m(\mathbf{y})$  designate an arbitrary probability tilting measure (which may depend on the transmitted codeword), then it follows that

$$\begin{aligned} P_{e|m} &\leq e^{nTs} \sum_{\mathbf{y}} \psi_n^m(\mathbf{y}) \psi_n^m(\mathbf{y})^{-1} p(\mathbf{y}|\mathbf{x}_m) \left( \sum_{m' \neq m} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^\rho \\ &\leq e^{nTs} \sum_{\mathbf{y}} \psi_n^m(\mathbf{y}) \left( \psi_n^m(\mathbf{y})^{-\frac{1}{\rho}} p(\mathbf{y}|\mathbf{x}_m)^{\frac{1}{\rho}} \sum_{m' \neq m} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^\rho. \end{aligned}$$

Next, invoking Jensen's inequality gives

$$P_{e|m} \leq e^{nTs} \left( \sum_{\mathbf{y}} \psi_n^m(\mathbf{y})^{1-\frac{1}{\rho}} p(\mathbf{y}|\mathbf{x}_m)^{\frac{1}{\rho}} \sum_{m' \neq m} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^\rho.$$

This concludes the proof of (9) by setting

$$\psi_n^m(\mathbf{y}) = \frac{G_n^m(\mathbf{y}) p(\mathbf{y}|\mathbf{x}_m)}{\sum_{\mathbf{y}} G_n^m(\mathbf{y}) p(\mathbf{y}|\mathbf{x}_m)} \quad (68)$$

where  $G_n^m(\mathbf{y})$  is an arbitrary non-negative function.

An undetected error event occurs if the received vector is included in the decision region of a codeword which differs from the transmitted codeword. Consequently, the average undetected error event satisfies

$$P_{ue} = \frac{1}{M} \sum_{m=1}^M \sum_{\mathbf{y} \in \Lambda_m} \sum_{m' \neq m} p(\mathbf{y}|\mathbf{x}_{m'}). \quad (69)$$

Note that in the case where list decoding is considered (i.e., the decision regions are not disjoint), the LHS of (69) is no longer a probability. However, for the latter case this expression equals the expected number of incorrect codewords in the decoded list. It follows from (69) that for  $0 \leq s \leq 1$ , the undetected error probability satisfies

$$\begin{aligned} P_{ue} &= \frac{1}{M} \sum_{m=1}^M \sum_{\mathbf{y} \in \Lambda_m} p(\mathbf{y}|\mathbf{x}_m) \left( \frac{\sum_{m' \neq m} p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^s \left( \frac{\sum_{m' \neq m} p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{1-s} \\ &\leq e^{nT(s-1)} \frac{1}{M} \sum_{m=1}^M \sum_{\mathbf{y}} p(\mathbf{y}|\mathbf{x}_m) \left( \sum_{m' \neq m} \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^s \end{aligned} \quad (70)$$

where the last inequality holds since for  $\mathbf{y} \in \Lambda_m$  and  $0 \leq s \leq 1$

$$\left( \frac{p(\mathbf{y}|\mathbf{x}_m)}{\sum_{m' \neq m} p(\mathbf{y}|\mathbf{x}_{m'})} \right)^{1-s} \geq e^{nT(1-s)}.$$

The rest of the proof follows in a similar way to the derivation of (9) when comparing the bound in (67) with (70).

APPENDIX E  
PROOF OF COROLLARY 1

Consider the ensemble of fully random block codes of length  $n$  symbols where the  $M = e^{nR}$  codewords of a codebook are chosen independently at random according to the probability distribution  $q_{\mathbf{X}}$  on  $\mathcal{X}^n$ .

Let  $D_{\{\mathbf{x}_i\}_{i=1}^M}(m, G_n^m, s, \rho)$  denote the functional  $D_{\mathbf{B}}(m, G_n^m, s, \rho)$  in (11) where the dependence on a specific codebook  $\{\mathbf{x}_i\}_{i=1}^M$  is expressed explicitly. Given a fixed codeword  $\mathbf{x}_m$  for the  $m$ -th message, the expectation over the other  $M - 1$  codewords on the right-hand side of (9) gives that for  $0 \leq s \leq \rho \leq 1$

$$\begin{aligned}
& \sum_{\{\mathbf{x}_i\}_{i=1}^M \setminus \{\mathbf{x}_m\}} \left( \prod_{i \neq m} q_{\mathbf{X}}(\mathbf{x}_i) \right) D_{\{\mathbf{x}_i\}_{i=1}^M}(m, G_n^m, s, \rho) \\
& \stackrel{(a)}{\leq} \left( \sum_{\mathbf{y}} G_n^m(\mathbf{y}) p(\mathbf{y}|\mathbf{x}_m) \right)^{1-\rho} \\
& \quad \left( \sum_{m' \neq m} \sum_{\mathbf{x}_{m'}} q_{\mathbf{X}}(\mathbf{x}_{m'}) \sum_{\mathbf{y}} p(\mathbf{y}|\mathbf{x}_m) G_n^m(\mathbf{y})^{1-\frac{1}{\rho}} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^{\rho} \\
& = (M-1)^{\rho} \left( \sum_{\mathbf{y}} G_n^m(\mathbf{y}) p(\mathbf{y}|\mathbf{x}_m) \right)^{1-\rho} \\
& \quad \left( \sum_{\mathbf{x}'} q_{\mathbf{X}}(\mathbf{x}') \sum_{\mathbf{y}} p(\mathbf{y}|\mathbf{x}_m) G_n^m(\mathbf{y})^{1-\frac{1}{\rho}} \left( \frac{p(\mathbf{y}|\mathbf{x}')}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^{\rho} \tag{71}
\end{aligned}$$

where (a) follows from (11) and by invoking Jensen's inequality. Next, by substituting the non-negative function

$$G_n^m(\mathbf{y}) \triangleq \left( \sum_{\mathbf{x}} q_{\mathbf{X}}(\mathbf{x}) \left( \frac{p(\mathbf{y}|\mathbf{x})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^{\rho}$$

in (71), one obtains that for  $0 \leq s \leq \rho \leq 1$  and  $m = 1, \dots, M$

$$\begin{aligned}
& \sum_{\{\mathbf{x}_i\}_{i=1}^M \setminus \{\mathbf{x}_m\}} \left( \prod_{i \neq m} q_{\mathbf{X}}(\mathbf{x}_i) \right) D_{\{\mathbf{x}_i\}_{i=1}^M}(m, G_n^m, s, \rho) \\
& \leq (M-1)^{\rho} \sum_{\mathbf{y}} p(\mathbf{y}|\mathbf{x}_m) \left( \sum_{\mathbf{x}'} q_{\mathbf{X}}(\mathbf{x}') \left( \frac{p(\mathbf{y}|\mathbf{x}')}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^{\rho}.
\end{aligned}$$

By averaging  $D_{\{\mathbf{x}_i\}_{i=1}^M}(m, G_n^m, s, \rho)$  over the  $M$  codewords, we get that for every index  $m$  ( $1 \leq m \leq M$ )

$$\begin{aligned}
& \sum_{\{\mathbf{x}_i\}_{i=1}^M} \left( \prod_{i=1}^M q_{\mathbf{X}}(\mathbf{x}_i) \right) D_{\{\mathbf{x}_i\}_{i=1}^M}(m, G_n^m, s, \rho) \\
& = \sum_{\mathbf{x}_m} q_{\mathbf{X}}(\mathbf{x}_m) \sum_{\{\mathbf{x}_i\}_{i=1}^M \setminus \{\mathbf{x}_m\}} \left( \prod_{i \neq m} q_{\mathbf{X}}(\mathbf{x}_i) \right) D_{\{\mathbf{x}_i\}_{i=1}^M}(m, G_n^m, s, \rho) \\
& \leq (M-1)^{\rho} \sum_{\mathbf{y}} \sum_{\mathbf{x}_m} q_{\mathbf{X}}(\mathbf{x}_m) p(\mathbf{y}|\mathbf{x}_m) \left( \sum_{\mathbf{x}'} q_{\mathbf{X}}(\mathbf{x}') \left( \frac{p(\mathbf{y}|\mathbf{x}')}{p(\mathbf{y}|\mathbf{x}_m)} \right)^{\frac{s}{\rho}} \right)^{\rho} \\
& = (M-1)^{\rho} \sum_{\mathbf{y}} \left\{ \left( \sum_{\mathbf{x}} q_{\mathbf{X}}(\mathbf{x}) p(\mathbf{y}|\mathbf{x})^{1-s} \right) \left( \sum_{\mathbf{x}'} q_{\mathbf{X}}(\mathbf{x}') p(\mathbf{y}|\mathbf{x}')^{\frac{s}{\rho}} \right)^{\rho} \right\}. \tag{72}
\end{aligned}$$

Since the right-hand side of (72) does not depend on the index  $m$ , then this bound also applies to the expectation of the quantity  $\frac{1}{M} \sum_{m=1}^M D_{\{\mathbf{x}_i\}_{i=1}^M}(m, G_n^m, s, \rho)$ . Therefore, there exists a block code for which the value of this

quantity is not larger than the average over the considered ensemble, i.e.,

$$\begin{aligned} & \frac{1}{M} \sum_{m=1}^M D_{\{\mathbf{x}_i\}_{i=1}^M} (m, G_n^m, s, \rho) \\ & \leq (M-1)^\rho \sum_{\mathbf{y}} \left\{ \left( \sum_{\mathbf{x}} q_{\mathbf{X}}(\mathbf{x}) p(\mathbf{y}|\mathbf{x})^{1-s} \right) \left( \sum_{\mathbf{x}'} q_{\mathbf{X}}(\mathbf{x}') p(\mathbf{y}|\mathbf{x}')^{\frac{s}{\rho}} \right)^\rho \right\}. \end{aligned} \quad (73)$$

From (9), (10) and (73), it follows that the above block code satisfies simultaneously

$$\begin{aligned} P_e &= \frac{1}{M} \sum_{m=1}^M P_{e|m} \\ &\leq e^{nsT} \cdot \frac{1}{M} \sum_{m=1}^M D_{\{\mathbf{x}_i\}_{i=1}^M} (m, G_n^m, s, \rho) \\ &\leq e^{nsT} (M-1)^\rho \sum_{\mathbf{y}} \left\{ \left( \sum_{\mathbf{x}} q_{\mathbf{X}}(\mathbf{x}) p(\mathbf{y}|\mathbf{x})^{1-s} \right) \left( \sum_{\mathbf{x}'} q_{\mathbf{X}}(\mathbf{x}') p(\mathbf{y}|\mathbf{x}')^{\frac{s}{\rho}} \right)^\rho \right\} \\ &< e^{n(sT+\rho R)} \sum_{\mathbf{y}} \left\{ \left( \sum_{\mathbf{x}} q_{\mathbf{X}}(\mathbf{x}) p(\mathbf{y}|\mathbf{x})^{1-s} \right) \left( \sum_{\mathbf{x}'} q_{\mathbf{X}}(\mathbf{x}') p(\mathbf{y}|\mathbf{x}')^{\frac{s}{\rho}} \right)^\rho \right\} \\ &= e^{-n(E_0(s,\rho,q_X) - \rho R - sT)} \end{aligned}$$

and

$$\begin{aligned} P_{ue} &< e^{n((s-1)T+\rho R)} \sum_{\mathbf{y}} \left\{ \left( \sum_{\mathbf{x}} q_{\mathbf{X}}(\mathbf{x}) p(\mathbf{y}|\mathbf{x})^{1-s} \right) \left( \sum_{\mathbf{x}'} q_{\mathbf{X}}(\mathbf{x}') p(\mathbf{y}|\mathbf{x}')^{\frac{s}{\rho}} \right)^\rho \right\} \\ &= e^{-n(E_0(s,\rho,q_X) - \rho R - (s-1)T)} \end{aligned}$$

where the last two equalities follow from (15), and since the input distribution and the channel are assumed to be memoryless, i.e.,

$$p(\mathbf{y}|\mathbf{x}) = \prod_{i=1}^n p(y_i|x_i), \quad q_{\mathbf{X}}(\mathbf{x}) = \prod_{i=1}^n q_X(x_i).$$

The proof of Corollary 1 is completed by optimizing the bounds over the parameters  $\rho$  and  $s$  (where  $0 \leq s \leq \rho \leq 1$ ) and the input distribution  $q_X$ . This gives the exponents  $E_1$  and  $E_2$  in (14) for the upper bounds on  $P_e$  and  $P_{ue}$ , respectively.

#### APPENDIX F PROOF OF PROPOSITION 6

The bounds in Proposition 6 are derived from Proposition 5 as follows: setting

$$p(\mathbf{y}|\mathbf{x}) = \prod_{i=1}^n p(y_i|x_i)$$

and

$$G_n^m(\mathbf{y}) = \prod_{i=1}^n g(y_i)$$

in (11), and relying on the useful rule for interchanging sum and product signs  $\sum_{\mathbf{y}} \prod_{i=1}^n f(y_i) = \prod_{i=1}^n \sum_{y_i} f(y_i)$ , one gets from (9) the RHS of (16) as an upper bound on  $P_{e|0}$ . Since the considered block code is linear and the communication channel is memoryless and symmetric, the bound in (16) follows from the message independence property in Proposition 2. The derivation of the bound in (17) relies on (10) where it is first proved that for a linear block code whose transmission takes place over a memoryless symmetric channel, the resulting expression

for  $D_B(m, G_n^m, s, \rho)$  is independent of  $m$ . To this end, let  $\mathcal{T}$  be a mapping as defined in Definition 1, then for all  $1 \leq i \leq n$

$$\begin{aligned} & \sum_{m' \neq m} \sum_{y \in \mathcal{Y}} g(y)^{1-\frac{1}{\rho}} p(y|x_{m,i}) \left( \frac{p(y|x_{m',i})}{p(y|x_{m,i})} \right)^{\frac{s}{\rho}} \\ &= \sum_{m' \neq m} \sum_{y \in \mathcal{Y}} g(y)^{1-\frac{1}{\rho}} p(\mathcal{T}(y, -x_{m,i})|0) \left( \frac{p(\mathcal{T}(y, -x_{m,i})|x_{m',i} - x_{m,i})}{p(\mathcal{T}(y, -x_{m,i})|0)} \right)^{\frac{s}{\rho}} \\ &= \sum_{l \neq 0} \sum_{z \in \mathcal{Y}} g(y)^{1-\frac{1}{\rho}} p(z|0) \left( \frac{p(z|x_{l,i})}{p(z|0)} \right)^{\frac{s}{\rho}}. \end{aligned}$$

As a result, it follows that for a memoryless and symmetric channel

$$\frac{1}{M} \sum_{m=1}^M D_B(m, G_n^m, s, \rho) = D(g, s, \rho) \quad (74)$$

where  $D(g, s, \rho)$  is introduced in (18). The proof of the upper bound on  $P_{ue}$  as given in (17) is completed by substituting (74) in (17).

## APPENDIX G

### ERROR AND RATE PERFORMANCE OF HYBRID-ARQ SCHEMES BASED ON GENERALIZED DECODING WITH ERASURES

The following discussion is provided in [15] and [19], and it is surveyed here for the sake of completeness.

Since Forney's generalized decoding rule (5) with a positive value of  $T$  is used in the context of erasures, the resulting decision regions at the receiver are disjoint, and the erasure probability  $P_x$  for a single block transmission is given by

$$P_x = P_e - P_{ue}$$

where  $P_e$  and  $P_{ue}$  are, respectively, the (total) block error probability and undetected error probability for a single block transmission. The erasure probability is studied via an upper bound on the error probability  $P_e$ . Assuming a noiseless and immediate feedback, for the case where no deadlines are considered, the expected rate of the considered system equals

$$(1 - P_x)R \quad (75)$$

where  $R$  is the rate of the codebook used (in units of bits per channel use) for a single block transmission. The error probability of this scheme is given by

$$\frac{P_{ue}}{1 - P_x}. \quad (76)$$

Note that the replacement of  $P_x$  in (75) and (76) with an upper bound on  $P_e$ , provides a lower bound on the expected rate and an upper bound on the error probability.

For the case where deadlines are considered, let  $Q$  ( $Q \geq 1$ ) be the maximal number of block retransmissions (including the first transmitted block). Each transmitted block is decoded separately using Forney's generalized decoding rule with erasures. Such a scheme is termed memoryless in [19] (note that the ARQ scheme without deadlines, studied in [15], is also memoryless in this sense). In cases where  $Q$  consequent block transmissions occur, then the generalized decoding rule is replaced for the last ( $Q$ -th) retransmitted block with an ML decoder. As a result, the expected rate and error probability, denoted by  $R(Q)$  and  $P_e(Q)$ , respectively, satisfy

$$\begin{aligned} R(Q) &= \frac{R}{\sum_{k=0}^{Q-1} (P_x)^k} \\ &= \frac{R(1 - P_x)}{1 - (P_x)^Q} \end{aligned} \quad (77)$$

and

$$\begin{aligned} P_e(Q) &= \sum_{k=1}^{Q-1} (P_x)^{k-1} P_{ue} + (P_x)^{Q-1} P_e^{\text{ML}} \\ &= \frac{(1 - (P_x)^{Q-1}) P_{ue}}{1 - P_x} + (P_x)^{Q-1} P_e^{\text{ML}} \end{aligned} \quad (78)$$

where  $P_e^{\text{ML}}$  is the block error probability under ML decoding for the considered code (while referring to the decoding of the last retransmitted block separately). Note that in the limit where  $Q \rightarrow \infty$  (no deadlines), then (77) and (78) tend asymptotically to (75) and (76), respectively. Replacing  $P_x$  in (77) and (78) with an upper bound on the (total) error probability  $P_e$ , results in a lower bound on the expected rate, and an upper bound on the error probability, respectively.

#### APPENDIX H PROOF OF PROPOSITION 7

*Proof of the upper bound on the conditional error probability in (35)*

Let  $\Lambda_m^{\text{LR}}$  designate the decision region in (7), then the conditional error probability is equal to

$$P_{e|m} = \sum_{\mathbf{y} \notin \Lambda_m^{\text{LR}}} p(\mathbf{y}|\mathbf{x}_m).$$

For  $\mathbf{y} \notin \Lambda_m^{\text{LR}}$ , the decision rule in (7) implies that

$$\frac{p(\mathbf{y}|\mathbf{x}_m)}{p(\mathbf{y}|\mathbf{x}_{m_2})} < e^{nT}$$

where  $\mathbf{x}_{m_2}$  is the second most probable codeword, and therefore

$$e^{nT} \sum_{m' \neq m} \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} > 1.$$

Let  $s \geq 0$ , then for  $\mathbf{y} \notin \Lambda_m^{\text{LR}}$

$$e^{nTs} \left( \sum_{m' \neq m} \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^s > 1$$

and the conditional block error probability satisfies

$$P_{e|m} \leq e^{nsT} \sum_{\mathbf{y}} p(\mathbf{y}|\mathbf{x}_m) \left( \sum_{m' \neq m} \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^s. \quad (79)$$

The bound in (35) follows from (79), using the arguments following (67).

*Proof of the upper bound on the conditional undetected error probability in (36)*

The conditional undetected error probability is given by

$$P_{ue|m} = \sum_{\mathbf{y} \in \mathcal{L}} p(\mathbf{y}|\mathbf{x}_m)$$

where

$$\mathcal{L} \triangleq \{ \mathbf{y} : \exists m' \neq m, p(\mathbf{y}|\mathbf{x}_{m'}) \geq e^{nT} p(\mathbf{y}|\mathbf{x}_{m_2}) \}$$

and  $\mathbf{x}_{m_2}$  is the second most probable codeword for  $p(\mathbf{y}|\mathbf{x})$ . Since  $p(\mathbf{y}|\mathbf{x}_{m_2}) \geq p(\mathbf{y}|\mathbf{x}_m)$ , then

$$\mathcal{L} \subseteq \{ \mathbf{y} : \exists m' \neq m, p(\mathbf{y}|\mathbf{x}_{m'}) \geq e^{nT} p(\mathbf{y}|\mathbf{x}_m) \}$$

and therefore

$$\begin{aligned}
\mathbf{y} \in \mathcal{L} &\Rightarrow \exists m' \neq m, \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \cdot e^{-nT} \geq 1 \\
&\Rightarrow e^{-nT} \sum_{m' \neq m} \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \geq 1 \\
&\Rightarrow \forall s \geq 0, e^{-nTs} \left( \sum_{m' \neq m} \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^s \geq 1
\end{aligned}$$

As a result, the conditional undetected block error probability satisfies, for all  $s \geq 0$ , the following upper bound:

$$P_{ue|m} \leq e^{-nsT} \sum_{\mathbf{y}} p(\mathbf{y}|\mathbf{x}_m) \left( \sum_{m' \neq m} \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^s.$$

The rest of the proof of (36) is, again, similar to the derivation following (67).

#### APPENDIX I PROOF OF PROPOSITION 8

The derivation of the bounds in Proposition 8 is primarily identical to the analysis in [9] and [26, Section 3.2.4], for which the reader is referred for a complete treatment of the analysis under ML decoding. We assume a BPSK modulation over AWGN channel with energy  $E_s$  per transmitted coded symbol, and a white Gaussian noise with two-sided power spectral density of  $\frac{N_0}{2}$ . Hence, the received vector  $\mathbf{y}$  satisfies

$$\mathbf{y} = \gamma \mathbf{x} + \mathbf{n} \quad (80)$$

where  $\gamma \triangleq \sqrt{\frac{2E_s}{N_0}}$ ,  $\mathbf{x} \in \mathcal{C} \subseteq \{-1, +1\}^n$  is the transmitted codeword (with BPSK modulation), and  $\mathbf{n}$  is a normal random vector with independent coordinates (all with zero mean and unit variance). Setting

$$E_e(d) \triangleq \left\{ \mathbf{y} \in \mathcal{Y}^n : \frac{\max_{\mathbf{x} \in \mathcal{C}_d \setminus \{\mathbf{x}_0\}} p(\mathbf{y}|\mathbf{x})}{p(\mathbf{y}|\mathbf{x}_0)} \cdot e^{nT} \geq 1 \right\}.$$

where  $\mathcal{C}_d$  is the set of all codewords whose Hamming weight is  $d$ , and  $\mathbf{x}_0$  is the all-zero codeword, it follows from (7) and the union bound that the conditional decoding error probability is upper bounded by

$$P_{e|0} \leq \sum_{d=d_{\min}}^n \Pr(E_e(d)) \quad (81)$$

where  $d_{\min}$  denotes the minimal Hamming distance of  $\mathcal{C}$ . Consider the following inequality on the probability of an error event:

$$\Pr(E) \leq \Pr(E, \mathbf{y} \in \mathcal{R}) + \Pr(\mathbf{y} \notin \mathcal{R}) \quad (82)$$

where  $E$  denotes an error event,  $\mathbf{y} \in \mathcal{Y}^n$  is the received vector, and  $\mathcal{R} \subseteq \mathbf{Y}^n$ . From (81) and (82), it follows that

$$P_{e|0} \leq \sum_{d=d_{\min}}^n \left( \Pr(E_e(d), \mathbf{y} \in \mathcal{R}) + \Pr(\mathbf{y} \notin \mathcal{R}) \right). \quad (83)$$

Using the union bound, we have

$$\begin{aligned}
\Pr(E_e(d), \mathbf{y} \in \mathcal{R}) &\leq \sum_{\mathbf{x} \in \mathcal{C}_d} \Pr \left( \frac{p(\mathbf{y}|\mathbf{x})}{p(\mathbf{y}|\mathbf{x}_0)} e^{nT} \geq 1, \mathbf{y} \in \mathcal{R} \right) \\
&\stackrel{(a)}{=} \sum_{\mathbf{x} \in \mathcal{C}_d} \Pr \left( \langle \mathbf{y}, \mathbf{x} \rangle \geq \langle \mathbf{y}, \mathbf{x}_0 \rangle - \frac{nT}{\gamma}, \mathbf{y} \in \mathcal{R} \right)
\end{aligned} \quad (84)$$

where equality (a) follows from (80), and  $\langle \mathbf{x}, \mathbf{y} \rangle \triangleq \sum_{i=1}^n x_i y_i$  denotes the scalar multiplication of the vectors  $\mathbf{x}$  and  $\mathbf{y}$ . Similarly to the derivation of bound in [9] (under ML decoding), we choose

$$\mathcal{R} \triangleq \left\{ \mathbf{y} : \|\mathbf{y} - \eta\gamma\mathbf{x}_0\|^2 \leq nr^2 \right\} \quad (85)$$

where  $\eta$  and  $r$  are arbitrary parameters which are subject to optimization. In addition, define

$$\begin{aligned} Z &\triangleq \langle \mathbf{y}, \mathbf{x} \rangle - \langle \mathbf{y}, \mathbf{x}_0 \rangle \\ W &\triangleq \|\mathbf{y} - \eta\gamma\mathbf{x}_0\|^2 - nr^2 \end{aligned}$$

then it follows from (84) and (85), using the Chernoff bound that

$$\Pr(E_e(d), \mathbf{y} \in \mathcal{R}) + \Pr(\mathbf{y} \notin \mathcal{R}) \leq e^{\frac{tnT}{\gamma}} |\mathcal{C}_d| \mathbb{E} \left[ e^{tZ+uW} \right] + \mathbb{E} \left[ e^{sW} \right] \quad (86)$$

for all  $t \geq 0$ ,  $u \leq 0$ , and  $s \geq 0$ . Evaluating the expectations in (86) and setting  $t = \frac{\gamma}{2}(1 - 2u\eta)$ , we have similarly to [9] and [26, Section 3.2.4]:

$$\begin{aligned} \Pr(E_e(d), \mathbf{y} \in \mathcal{R}) + \Pr(\mathbf{y} \notin \mathcal{R}) &\leq e^{\frac{nT(1-r\eta)}{2}} |\mathcal{C}_d| e^{-nur^2} (f_1(\gamma, u, \eta))^{n-d} (f_2(\gamma, u, \eta))^d \\ &\quad + e^{-nsr^2} (f_1(\gamma, s, \eta))^n \end{aligned} \quad (87)$$

where

$$\begin{aligned} f_1(\gamma, \alpha, \eta) &\triangleq \frac{e^{\frac{(1-\eta)^2 \gamma^2 \alpha}{1-2\alpha}}}{\sqrt{1-2\alpha}} \\ f_2(\gamma, \alpha, \eta) &\triangleq \frac{e^{-\frac{\gamma^2(1-2\alpha\eta^2)}{2}}}{\sqrt{1-2\alpha}}, \quad \alpha < \frac{1}{2}. \end{aligned}$$

Optimizing the term  $e^{nr^2}$  on the right-hand side of (87), gives

$$\Pr(E_e(d), \mathbf{y} \in \mathcal{R}) + \Pr(\mathbf{y} \notin \mathcal{R}) \leq 2^{h_2\left(\frac{s}{s-u}\right)} A^{-\frac{u}{s-u}} B^{-\frac{s}{s-u}}, \quad 0 < s < \frac{1}{2}, \quad u \leq 0 \quad (88)$$

where

$$\begin{aligned} A &\triangleq (f_1(\gamma, s, \eta))^n \\ B &\triangleq e^{\frac{nT(1-r\eta)}{2}} |\mathcal{C}_d| (f_1(\gamma, u, \eta))^{n-d} (f_2(\gamma, u, \eta))^d \end{aligned}$$

and  $h_2$  designates the binary entropy function on base 2. Using the change of variables

$$\begin{aligned} \rho &\triangleq \frac{s}{s-u} \\ \beta &\triangleq \rho(1-2u) \\ \xi &\triangleq \rho(1-2u\eta) \end{aligned}$$

where  $0 \leq \rho \leq 1$ ,  $0 \leq \beta \leq 1$ , and  $\xi \geq 0$ , the bound in (88) transforms to

$$\Pr(E_e(d), \mathbf{y} \in \mathcal{R}) + \Pr(\mathbf{y} \notin \mathcal{R}) \leq 2^{h_2(\rho)} e^{-nE(E_s/N_0, d/n, \beta, \rho, \xi) + \frac{nT\xi}{2}} \quad (89)$$

where

$$E(c, \delta, \beta, \rho, \xi) \triangleq -\rho r_n(\delta) - \frac{\rho}{2} \ln \left( \frac{\rho}{\beta} \right) - \frac{1-\rho}{2} \ln \left( \frac{1-\rho}{1-\beta} \right) + c \left( 1 - (1-\delta) \frac{\xi^2}{\beta} - \frac{(1-\xi)^2}{1-\beta} \right).$$

The parameters  $\rho$ ,  $\beta$  and  $\xi$  are optimized in [9], [26] such that the error exponent  $E(c, \delta, \beta, \rho, \xi)$  is maximized<sup>3</sup> (note that the bound for  $T = 0$  coincides with the bound which refers to ML decoding), setting the optimal parameters

<sup>3</sup>It is possible to obtain the optimized  $\rho$  and  $\xi$  when maximizing the entire exponent  $E(c, \delta, \beta, \rho, \xi) + \frac{T\xi}{2}$ . To this end,  $\xi$  needs to be shifted by  $-\frac{T}{2}$  and the optimal  $\rho$  remains without change. The parameter  $\beta$  is required to be numerically optimized over  $0 \leq \beta \leq 1$ . Nevertheless, the resulting bound gives only a marginal gain over the bound which maximizes  $E(c, \delta, \beta, \rho, \xi)$  without the addition of  $\frac{T\xi}{2}$ .



yields the first argument in (41). The second term inside the minimization on the right-hand side of (41) follows from a union bound on the error probability

$$P_e \leq \sum_{d=d_{\min}}^n \sum_{\mathbf{x} \in \mathcal{C}_d} \Pr \left( \frac{p(\mathbf{y}|\mathbf{x})}{p(\mathbf{y}|\mathbf{x}_0)} e^{nT} \geq 1 \right)$$

where for every codeword  $\mathbf{x} \in \mathcal{C}_d$

$$\Pr \left( \frac{p(\mathbf{y}|\mathbf{x})}{p(\mathbf{y}|\mathbf{x}_0)} e^{nT} \geq 1 \right) = Q \left( \gamma\sqrt{d} - \frac{nT}{2\gamma\sqrt{d}} \right).$$

The derivation of the upper bound on the undetected error probability follows some similar arguments, and is therefore omitted.

## APPENDIX J PROOF OF PROPOSITION 9

The main ingredient for proving the DS2 bound on the block error probability under ML decoding (and also the well known random-coding bound) is that for a received vector  $\mathbf{y}$  which is not included in the decision region  $\Lambda_m$  as given in (3), the following inequality holds:

$$1 \leq \left( \sum_{m' \neq m} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^\lambda \right)^\rho, \quad \lambda, \rho \geq 0. \quad (90)$$

When an error event under fixed-size ( $L$ ) list decoding is considered, there exists  $L$  distinct codewords, all different from the transmitted codeword, whose a-posteriori probability is larger than the one of the transmitted codeword. Hence, the sum on the right-hand side of (90) is divided by  $L$ . Specifically for a received vector  $\mathbf{y}$  that results in an error event, the following inequality is satisfied:

$$1 \leq \left( \frac{1}{L} \sum_{m' \neq m} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^\lambda \right)^\rho, \quad \lambda, \rho \geq 0 \quad (91)$$

Following the derivation of the DS2 bound in [26, p. 96] where the right-hand side of (90) is replaced with (91) leads to the derivation of the bound in Proposition 9. This derivation is repeated for the sake of completeness. For an arbitrarily chosen probability measure  $\psi_n^m(\mathbf{y})$  it follows that:

$$\begin{aligned} P_{e|m} &\leq \sum_{\mathbf{y}} \psi_n^m(\mathbf{y}) (\psi_n^m(\mathbf{y}))^{-1} p(\mathbf{y}|\mathbf{x}_m) \left( \frac{1}{L} \sum_{m' \neq m} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^\lambda \right)^\rho \\ &= \sum_{\mathbf{y}} \psi_n^m(\mathbf{y}) \left( (\psi_n^m(\mathbf{y}))^{-\frac{1}{\rho}} (p(\mathbf{y}|\mathbf{x}_m))^{\frac{1}{\rho}} \frac{1}{L} \sum_{m' \neq m} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^\lambda \right)^\rho \\ &\leq \left( \sum_{m' \neq m} \sum_{\mathbf{y}} (\psi_n^m(\mathbf{y}))^{1-\frac{1}{\rho}} (p(\mathbf{y}|\mathbf{x}_m))^{\frac{1}{\rho}} \frac{1}{L} \left( \frac{p(\mathbf{y}|\mathbf{x}_{m'})}{p(\mathbf{y}|\mathbf{x}_m)} \right)^\lambda \right)^\rho \end{aligned}$$

where the last inequality follows from Jensen's inequality. Plugging  $\psi_n^m(\mathbf{y})$  as in (68) concludes the proof.

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