

Upper Bounds on the Relative Entropy and Rényi Divergence as a Function of Total Variation Distance for Finite Alphabets

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Cast of Characters

- A finite set \mathcal{A} with σ -algebra \mathcal{F}
- Probability measures P, Q defined on the measurable space $(\mathcal{A}, \mathcal{F})$.
- $X \sim P$.
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Total Variation (TV) Distance

$$\begin{aligned} |P - Q| &= 2 \sup_{\mathcal{F} \in \mathcal{F}} (P(\mathcal{F}) - Q(\mathcal{F})) \\ &= \sum_{x \in \mathcal{A}} |P(x) - Q(x)|. \end{aligned}$$

Csiszár-Kemperman-Kullback-Pinsker Inequality

$$D(P\|Q) \geq \frac{1}{2} |P - Q|^2 \log e$$

and the constant is tight in the sense that

$$\inf_{P \neq Q} \frac{D(P\|Q)}{|P - Q|^2} = \frac{1}{2} \log e.$$

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An Implication of Pinsker's Inequality

Convergence in relative entropy \implies convergence in TV distance.

Question

Is there a reverse Pinsker inequality that provides an upper bound on the relative entropy as a function of the TV distance ?

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No, for every $\varepsilon > 0$ there exist P, Q s.t. $|P - Q| \leq \varepsilon$, $D(P||Q) = \infty$.

A Reverse Pinsker Inequality [Csiszár and Talata, 2006]

If $Q_{\min} \triangleq \min_{x \in \mathcal{A}} Q(x) > 0$, then

$$D(P\|Q) \leq \frac{\log e}{Q_{\min}} \cdot |P - Q|^2.$$

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A Reverse Pinsker Inequality [Verdú, ITA 2014]

If $P \ll Q$, let $\beta_1 \in [0, 1]$ be given by $\beta_1^{-1} \triangleq \sup_{x \in \mathcal{A}} \frac{dP}{dQ}(x)$.

Then, if $\beta_1 < 1$ as above,

$$D(P\|Q) \leq \frac{\log \frac{1}{\beta_1}}{2(1 - \beta_1)} \cdot |P - Q|.$$

Hence, if the relative information $i_{P\|Q}(x) \triangleq \log \frac{dP}{dQ}(x)$ is bounded from above, a reverse Pinsker inequality exists.

New Reverse Pinsker Inequality

Theorem

a) If $P \ll Q$

$$D(P\|Q) \leq \log \left(1 + \frac{|P - Q|^2}{2Q_{\min}} \right). \quad (1)$$

b) Furthermore, if $Q \ll P$ and $\beta_2 \in [0, 1]$ is given by

$$\beta_2 = \min_{x \in \mathcal{A}} \frac{P(x)}{Q(x)}$$

then the following tightened bound holds:

$$D(P\|Q) \leq \log \left(1 + \frac{|P - Q|^2}{2Q_{\min}} \right) - \frac{\beta_2 \log e}{2} \cdot |P - Q|^2. \quad (2)$$

Note on the Weaker Version of This Bound

This already improves the Csiszár-Talata inequality since

$$\log \left(1 + \frac{|P - Q|^2}{2Q_{\min}} \right) \leq \frac{\log e}{2Q_{\min}} \cdot |P - Q|^2.$$

Proof

The idea is to obtain upper and lower bounds on the χ^2 -divergence

$$\chi^2(P, Q) \triangleq \sum_{x \in \mathcal{A}} \frac{(P(x) - Q(x))^2}{Q(x)}.$$

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Bounds on the χ^2 -divergence

$$\begin{aligned} \chi^2(P \| Q) &= e^{D_2(P \| Q)} - 1 \\ &\geq e^{D(P \| Q)} - 1. \end{aligned}$$

$$\begin{aligned} \chi^2(P \| Q) &\leq \frac{\sum_{x \in \mathcal{A}} (P(x) - Q(x))^2}{Q_{\min}} \\ &\leq \frac{|P - Q|^2}{2Q_{\min}}. \end{aligned}$$

Combining the bounds yields the weaker version of this inequality.

Distance From the Equiprobable Distribution

If P is a distribution on a finite set \mathcal{A} , $H(P)$ gauges the “distance” from U , the equiprobable distribution defined on \mathcal{A} , since

$$H(P) = \log |\mathcal{A}| - D(P\|U).$$

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Notation

- $h(x) = -x \log x - (1 - x) \log(1 - x)$ for $x \in [0, 1]$ denotes the binary entropy function;
- $d(x||y) = x \log \frac{x}{y} + (1 - x) \log \frac{1-x}{1-y}$ for $x, y \in [0, 1]$ denotes the binary relative entropy.

Theorem

Let U be the equiprobable distribution on a $\{1, \dots, |\mathcal{A}|\}$, with $1 < |\mathcal{A}| < \infty$.

a) For $\Delta \in (0, 2(1 - |\mathcal{A}|^{-1})]$,

$$\max_{P: |P-U|=\Delta} H(P) = \log |\mathcal{A}| - \min_m d(|\mathcal{A}|^{-1}m + \frac{1}{2}\Delta \parallel |\mathcal{A}|^{-1}m) \quad (3)$$

where the minimum in the right side of (3) is over all positive integers m not exceeding $\min\{|\mathcal{A}| - 1, |\mathcal{A}|(1 - \frac{1}{2}\Delta)\}$. Denoting such an integer by m_Δ , the maximum in the left side of (3) is attained by

$$P_\Delta(\ell) = \begin{cases} |\mathcal{A}|^{-1} + \frac{\Delta}{2m_\Delta}, & \ell \in \{1, \dots, m_\Delta\}, \\ |\mathcal{A}|^{-1} - \frac{\Delta}{2(|\mathcal{A}| - m_\Delta)}, & \ell \in \{m_\Delta + 1, \dots, |\mathcal{A}|\}. \end{cases} \quad (4)$$

Theorem (cont.)

b) Let

$$h_k = \begin{cases} 0, & k = 0 \\ h(|\mathcal{A}|^{-1}k) + |\mathcal{A}|^{-1}k \log k, & k \in \{1, \dots, |\mathcal{A}| - 2\} \\ \log |\mathcal{A}|, & k = |\mathcal{A}| - 1. \end{cases} \quad (5)$$

If $H \in [h_{k-1}, h_k)$ for $k \in \{1, \dots, |\mathcal{A}| - 1\}$, then

$$\min_{P: H(P)=H} |P - \mathbf{U}| = 2(1 - (k + \theta) |\mathcal{A}|^{-1}) \quad (6)$$

which is achieved by

$$P_\theta^{(k)}(\ell) = \begin{cases} 1 - (k - 1 + \theta) |\mathcal{A}|^{-1}, & \ell = 1 \\ |\mathcal{A}|^{-1}, & \ell \in \{2, \dots, k\}, \\ \theta |\mathcal{A}|^{-1}, & \ell = k + 1 \\ 0, & \ell \in \{k + 2, \dots, |\mathcal{A}|\} \end{cases} \quad (7)$$

where $\theta \in [0, 1)$ is chosen so that $H(P_\theta^{(k)}) = H$.

Bounds

Most well-known bound on the entropy difference in terms of the total variation distance is

$$|H(P) - H(Q)| \leq |P - Q| \log \left(\frac{|\mathcal{A}|}{|P - Q|} \right) \quad (8)$$

which holds if P, Q are probability measures defined on a finite set \mathcal{A} with $|P(a) - Q(a)| \leq \frac{1}{2}$ for all $a \in \mathcal{A}$.

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Particularizing (8) to the case where $Q = \mathbf{U}$, and $|P(a) - \frac{1}{|\mathcal{A}|}| \leq \frac{1}{2}$ for all $a \in \mathcal{A}$ yields

$$H(P) \geq \log |\mathcal{A}| - |P - \mathbf{U}| \log \left(\frac{|\mathcal{A}|}{|P - \mathbf{U}|} \right). \quad (9)$$

Bounds (cont.)

- The Bretagnolle-Huber inequality (1979)

$$\frac{1}{4}|P - Q|^2 \leq 1 - \exp(-D(P\|Q)). \quad (10)$$

- Particularizing Pinsker's inequality, (10) and the reverse Pinsker inequality yield

$$H(P) \leq \log |\mathcal{A}| - \frac{1}{2} |P - U|^2 \log e, \quad (11)$$

$$H(P) \leq \log |\mathcal{A}| + \log \left(1 - \frac{1}{4} |P - U|^2\right), \quad (12)$$

$$H(P) \geq \log |\mathcal{A}| - \log \left(1 + \frac{|\mathcal{A}|}{2} |P - U|^2\right). \quad (13)$$

- It can be checked that the lower bound on $H(P)$ in the previous slide is worse than (13), irrespectively of $|P - U|$, if either $|\mathcal{A}| = 2$ or $8 \leq |\mathcal{A}| \leq 102$ (note that $0 \leq |P - U| \leq 2(1 - |\mathcal{A}|^{-1})$).

Exact Locus versus Bounds

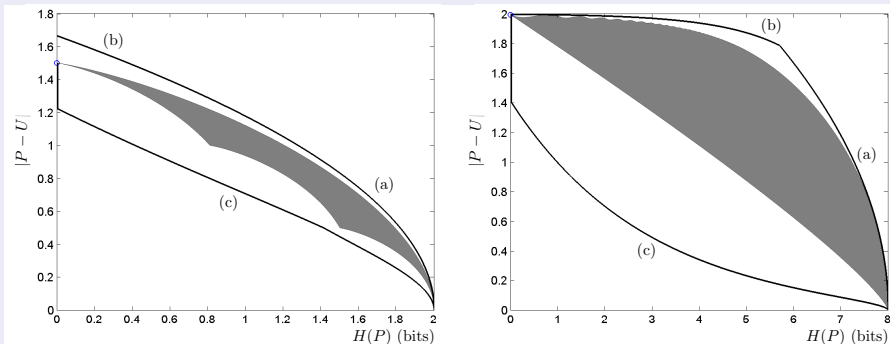


Figure: The exact locus of $(H(P), |P - U|)$ among all the probability measures P defined on a finite set \mathcal{A} , and bounds on $|P - U|$ as a function of $H(P)$ for $|\mathcal{A}| = 4$ (left plot), and $|\mathcal{A}| = 256$ (right plot). The point $(H(P), |P - U|) = (0, 2(1 - |\mathcal{A}|^{-1}))$ is depicted on the y -axis. In the two plots, Curves (a), (b) and (c) refer, respectively, to (11), (12) and (13); the exact locus (shaded region) refers to Theorem 2.

The Rényi Divergence

In the discrete case, we have for $\alpha \in (0, 1) \cup (1, \infty)$

$$D_\alpha(P||Q) = \frac{1}{\alpha - 1} \log \left(\sum_{x \in \mathcal{A}} P^\alpha(x) Q^{1-\alpha}(x) \right).$$

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Extreme cases:

- If $\alpha = 0$ then $D_0(P||Q) = -\log Q(\text{Support}(P))$,
- If $\alpha = +\infty$ then $D_\infty(P||Q) = \log \left(\text{ess sup } \frac{P}{Q} \right)$.
- If $D(P||Q) < \infty$, then $D(P||Q) = \lim_{\alpha \rightarrow 1^-} D_\alpha(P||Q)$.

New Bounds on the Rényi Divergence

Theorem

Let $P \ll\ll Q$ be defined a common finite set \mathcal{A} , and assume that P, Q are strictly positive with minimum masses denoted by P_{\min} and Q_{\min} , respectively. Let $\beta_1 = \min_{a \in \mathcal{A}} Q(a)/P(a)$ and $\beta_2 = \min_{a \in \mathcal{A}} P(a)/Q(a)$. Abbreviate $\delta \triangleq \frac{1}{2}|P - Q| \in [0, 1]$. Then, the Rényi divergence of order $\alpha \in [0, \infty]$ satisfies

$$D_\alpha(P\|Q) \leq \begin{cases} f_1, & \alpha \in (2, \infty] \\ f_2, & \alpha \in [1, 2] \\ \min \{f_2, f_3, f_4\}, & \alpha \in (\frac{1}{2}, 1) \\ \min \left\{ 2 \log \left(\frac{1}{1-\delta} \right), f_2, f_3, f_4 \right\}, & \alpha \in [0, \frac{1}{2}] \end{cases} \quad (14)$$

Theorem (cont.)

For $\alpha \in [0, \infty]$,

$$f_1(\alpha, \beta_1, \delta) \triangleq \begin{cases} \frac{1}{\alpha-1} \log \left(1 + \frac{\delta(\beta_1^{1-\alpha}-1)}{1-\beta_1} \right), & \alpha \in [0, 1) \cup (1, \infty) \\ \frac{\delta}{1-\beta_1} \log \frac{1}{\beta_1}, & \alpha = 1, \\ \log \frac{1}{\beta_1}, & \alpha = \infty \end{cases} \quad (15)$$

for $\alpha \in [0, 2]$

$$f_2(\alpha, \beta_1, Q_{\min}, \delta) \triangleq \min \left\{ f_1(\alpha, \beta_1, \delta), \log \left(1 + \frac{2\delta^2}{Q_{\min}} \right) \right\} \quad (16)$$

and, for $\alpha \in [0, 1)$, f_3 and f_4 are given by

$$f_3(\alpha, P_{\min}, \beta_1, \delta) \triangleq \left(\frac{\alpha}{1-\alpha} \right) \left[\log \left(1 + \frac{2\delta^2}{P_{\min}} \right) - 2\beta_1\delta^2 \log e \right], \quad (17)$$

$$f_4(\beta_2, Q_{\min}, \delta) \triangleq \log \left(1 + \frac{2\delta^2}{Q_{\min}} \right) - 2\beta_2\delta^2 \log e. \quad (18)$$

Relation with a Reverse Pinsker Inequality

By letting $\alpha \rightarrow 1$ in Theorem 3, we get

$$D(P\|Q) \leq \frac{\log \frac{1}{\beta_1}}{2(1 - \beta_1)} \cdot |P - Q|$$

which coincides with the bound by Verdú (ITA '14). The bound on the Rényi divergence in Theorem 3 therefore generalizes the above bound on the relative entropy.

Plot: Bounds on the Rényi Divergence

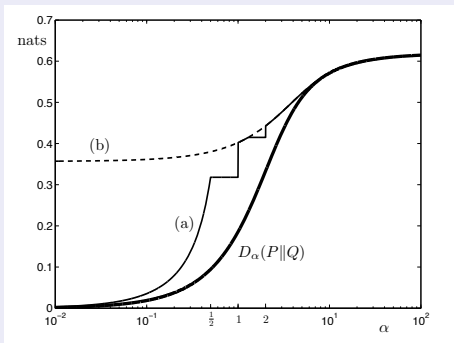


Figure: The Rényi divergence $D_\alpha(P\|Q)$ for P and Q which are defined on a binary alphabet with $P(0) = Q(1) = 0.65$, compared to (a) its upper bound in Theorem 3, and (b) its upper bound: $D_\alpha(P\|Q) \leq \frac{1}{\alpha-1} \log \left(1 + \frac{|P-Q|}{2} \frac{\beta_1^{1-\alpha}-1}{1-\beta_1} \right)$. These bounds coincide here when $\alpha \in (1, 1.291) \cup (2, \infty)$.

Summary

- This talk presents in part bounds among f -divergences which are derived for finite alphabets.
- The full paper version, which is currently under review and is available at <http://arxiv.org/abs/1508.00335>, derives bounds among f -divergences (and Rényi divergences) for general alphabets.