

Fast Polarization for Processes with Memory

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June 2018

In this talk

Setting: binary-input, symmetric, memoryless channel

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Polar codes: [Arıkan:09], [ArıkanTelatar:09], [Şaçoğlu+:09],

[KoradaUrbanke:10], [HondaYamamoto:13]

- ▶ **Setting:** **Memoryless i.i.d.** process $(X_i, Y_i)_{i=1}^N$
- ▶ **For simplicity:** Assume X_i binary
- ▶ **Polar transform:** $U_1^N = X_1^N \cdot G_N$
- ▶ **Index sets:**

$$\text{Low entropy: } \Lambda_N = \left\{ i : H(U_i | U_1^{i-1}, Y_1^N) < 2^{-N^\beta} \right\}$$

$$\text{High entropy: } \Omega_N = \left\{ i : H(U_i | U_1^{i-1}, Y_1^N) > 1 - 2^{-N^\beta} \right\}$$

- ▶ **Polarization:**

$$\lim_{N \rightarrow \infty} \frac{1}{N} |\Lambda_N| = 1 - H(X_1 | Y_1)$$

$$\lim_{N \rightarrow \infty} \frac{1}{N} |\Omega_N| = H(X_1 | Y_1)$$

Polar codes:

Optimal rate for:

- ▶ Coding for non-symmetric **memoryless** channels
- ▶ Coding for **memoryless** channels with non-binary inputs
- ▶ (Lossy) compression of **memoryless** sources

Question

- ▶ How to handle memory?

A framework for memory

- ▶ **Process:**

$$(X_i, Y_i, S_i)_{i=1}^N$$

- ▶ **Finite number of states:** $S_i \in \mathcal{S}$, where $|\mathcal{S}| < \infty$
- ▶ **Hidden state:** S_i is **unknown** to encoder and decoder

- ▶ **Probability distribution:**

$$P(x_i, y_i, s_i | s_{i-1})$$

- ▶ **Stationary:** same for all i
- ▶ **Markov:**

$$P(x_i, y_i, s_i | s_{i-1}) = P(x_i, y_i, s_i | \{x_j, y_j, s_j\}_{j < i})$$

- ▶ **State sequence:** aperiodic and irreducible Markov chain

Example 1

- ▶ **Model:** Finite state channel

$$P_s(y|x), \quad s \in \mathcal{S}$$

- ▶ **Input distribution:** X_j i.i.d. and independent of state
- ▶ **State transition:**

$$\pi(s_j | s_{j-1})$$

- ▶ **Distribution:**

$$P(x_i, y_i, s_j | s_{j-1}) = P(x_i) \pi(s_j | s_{j-1}) P_{s_j}(y_i | x_i)$$

Example 2

- ▶ **Model:** ISI + noise

$$Y_i = h_0 X_i + h_1 X_{i-1} + \dots + h_m X_{i-m} + \text{noise}$$

- ▶ **Input:** X_i has memory

$$P(X_i | X_{i-1}, X_{i-2}, \dots, X_{i-m}, X_{i-m-1})$$

- ▶ **State:**

$$S_i = [X_i \quad X_{i-1} \quad \dots \quad X_{i-m}]$$

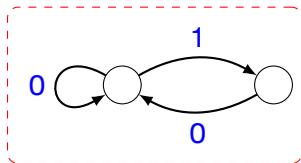
- ▶ **Distribution:** For X_i, S_i, S_{i-1} compatible,

$$P(X_i, Y_i, S_i | S_{i-1}) = P_{\text{noise}}(Y_i | h^T S_i) \cdot P(X_i | S_{i-1})$$

Example 3

- **Model:** (d, k) -RLL constrained system with noise

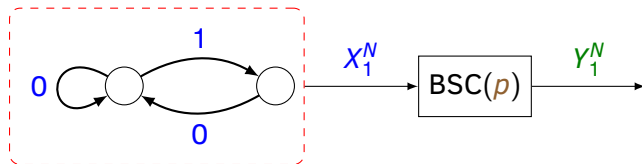
$(1, \infty)$ -RLL Constraint



Example 3

- **Model:** (d, k) -RLL constrained system with noise

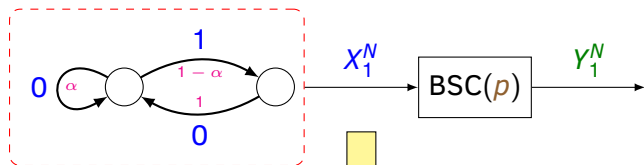
$(1, \infty)$ -RLL Constraint



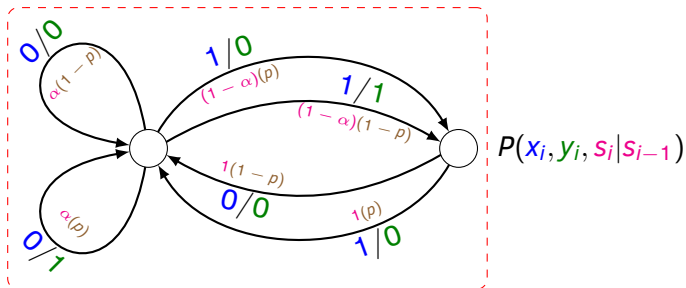
Example 3

- **Model:** (d, k) -RLL constrained system with noise

$(1, \infty)$ -RLL Constraint



state Markov chain



Polar codes: [Şaşoğlu:11], [ŞaşoğluTal:16], [ShuvalTal:17]

- ▶ **Setting:** Process $(X_i, Y_i, S_i)_{i=1}^N$ **with memory**, as above
- ▶ **Hidden state:** State **unknown** to encoder and decoder
- ▶ **Polar transform:** $U_1^N = X_1^N \cdot G_N$
 U_1^N are neither independent, nor identically distributed
- ▶ **Index sets:**


Low entropy: $\Lambda_N = \left\{ i : H(U_i | U_1^{i-1}, Y_1^N) < 2^{-N^\beta} \right\}$

High entropy: $\Omega_N = \left\{ i : H(U_i | U_1^{i-1}, Y_1^N) > 1 - 2^{-N^\beta} \right\}$

- ▶ **Polarization:**

$$\lim_{N \rightarrow \infty} \frac{1}{N} |\Lambda_N| = 1 - H_*(X|Y)$$

$$\lim_{N \rightarrow \infty} \frac{1}{N} |\Omega_N| = H_*(X|Y)$$


$$\lim_{N \rightarrow \infty} \frac{1}{N} H(X_1^N | Y_1^N)$$

Achievable rate

- ▶ **Achievable rate:** In all examples, R approaches

$$I_*(X; Y) = \lim_{N \rightarrow \infty} \frac{1}{N} I(X_1^N; Y_1^N)$$

- ▶ Also lossy compression of a source with memory
- ▶ **Successive cancellation:** [Wang+:15]
- ▶ Without state estimation!

Three parameters

- ▶ Joint distribution $P(x, y)$
- ▶ For simplicity: $X \in \{0, 1\}$
- ▶ **Parameters:**

Entropy $H(X|Y) = - \sum_{x,y} P(x, y) \log P(x|y)$

Bhattacharyya $Z(X|Y) = 2 \sum_y \sqrt{P(0, y)P(1, y)}$

T.V. distance $K(X|Y) = \sum_y |P(0, y) - P(1, y)|$

- ▶ **Connections:**

$$H \approx 0 \iff Z \approx 0 \iff K \approx 1$$

$$H \approx 1 \iff Z \approx 1 \iff K \approx 0$$

Three processes

For $n = 1, 2, \dots$

- ▶ $N = 2^n$
- ▶ $U_1^N = X_1^N G_N$
- ▶ Pick $B_n \in \{0, 1\}$ uniform, i.i.d.
- ▶ **Random index** from $\{1, 2, \dots, N\}$

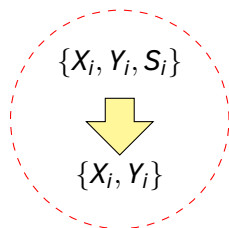
$$i = 1 + \langle B_1 B_2 \cdots B_n \rangle_2$$

- ▶ **Processes:**

Entropy $H_n = H(U_i | U_1^{i-1}, Y_1^N)$

Bhattacharyya $Z_n = Z(U_i | U_1^{i-1}, Y_1^N)$

T.V. distance $K_n = K(U_i | U_1^{i-1}, Y_1^N)$



Proof — memoryless case

Slow polarization

$$H_n \in (\epsilon, 1 - \epsilon)$$



$$|H_{n+1} - H_n| > 0$$



Fast polarization

$$Z_{n+1} \leq \begin{cases} 2Z_n & B_{n+1} = 0 \\ Z_n^2 & B_{n+1} = 1 \end{cases}$$



$$\frac{1}{N} |\Lambda_N| \xrightarrow{n \rightarrow \infty} 1 - H(X_1 | Y_1)$$

Low entropy set

New

$$K_{n+1} \leq \begin{cases} K_n^2 & B_{n+1} = 0 \\ 2K_n & B_{n+1} = 1 \end{cases}$$



$$\frac{1}{N} |\Omega_N| \xrightarrow{n \rightarrow \infty} H(X_1 | Y_1)$$

High entropy set

Proof — ~~memoryless~~ case [ŞaşoğluTal:16], [ShuvalTal:17]

Slow polarization

$$H_n \in (\epsilon, 1 - \epsilon)$$



$$|H_{n+1} - H_n| > 0$$

Fast polarization

$$Z_{n+1} \leq \begin{cases} 2\psi Z_n & B_{n+1} = 0 \\ \psi Z_n^2 & B_{n+1} = 1 \end{cases}$$



$$\frac{1}{N} |\Lambda_N| \xrightarrow{n \rightarrow \infty} 1 - H_*(X|Y)$$

Low entropy set

$$\psi = \psi(0) = \max_s \frac{1}{\pi(s)}$$

π : stationary state distribution

New

$$\hat{K}_{n+1} \leq \begin{cases} \psi \hat{K}_n^2 & B_{n+1} = 0 \\ 2\hat{K}_n & B_{n+1} = 1 \end{cases}$$



$$\frac{1}{N} |\Omega_N| \xrightarrow{n \rightarrow \infty} H_*(X|Y)$$

High entropy set

$\{X_i, Y_i, S_i\}$



$\{X_i, Y_i\}$

Fast polarization to high entropy set Ω_N

- ▶ **Memoryless case:**

- ▶ Parameter evolution inequality hinges on independence:

$$P(x_1^{2N}, y_1^{2N}) = P(x_1^N, y_1^N) \cdot P(x_{N+1}^{2N}, y_{N+1}^{2N})$$

- ▶ **Memory case:**

- ▶ **Force** independence: condition on middle state S_N

$$P(x_1^{2N}, y_1^{2N} | S_N) = P(x_1^N, y_1^N | S_N) \cdot P(x_{N+1}^{2N}, y_{N+1}^{2N} | S_N)$$

- ▶ **New processes:**

$$\hat{H}_n = H(U_i | U_1^{i-1}, Y_1^N, S_0, S_N)$$

$$\hat{K}_n = K(U_i | U_1^{i-1}, Y_1^N, S_0, S_N)$$

Polarization of K_n (memoryless case)

$$U_1^N = X_1^N \cdot G_N$$

$$V_1^N = X_{N+1}^{2N} \cdot G_N$$

$$Q_i = (U_1^{i-1}, Y_1^N)$$

$$R_i = (V_1^{i-1}, Y_{N+1}^{2N})$$

- ▶ **Memoryless** assumption:

$$P(u_i, v_i, q_i, r_i) = P(u_i, q_i) \cdot P(v_i, r_i)$$

- ▶ **Notation:**

$$T_i = U_i + V_i$$

- ▶ **One step polarization:**

$$K_{n+1} = \begin{cases} K(T_i | Q_i, R_i) & B_{n+1} = 0 & \text{'-' transform} \\ K(V_i | T_i, Q_i, R_i) & B_{n+1} = 1 & \text{'+' transform} \end{cases}$$

- ▶ **Recall:**

$$K(X|Y) = \sum_y |P(0,y) - P(1,y)|$$

Polarization of K_n (memoryless case), ‘-’ transform

$$\begin{aligned}K_{n+1} &= \sum_{q,r} |P_{T_i, Q_i, R_i}(0, q, r) - P_{T_i, Q_i, R_i}(1, q, r)| \\&= \sum_{q,r} \left| \sum_{v=0}^1 P(v, r)(P(v, q) - P(v+1, q)) \right| \\&= \sum_{q,r} \left| (P(0, q) - P(1, q))(P(0, r) - P(1, r)) \right| \\&= \sum_{q,r} |P(0, q) - P(1, q)| \cdot |P(0, r) - P(1, r)| \\&= \sum_q |P(0, q) - P(1, q)| \cdot \sum_r |P(0, r) - P(1, r)| \\&= K_n^2,\end{aligned}$$

Polarization of K_n (memoryless case), '+' transform

$$\begin{aligned}K_{n+1} &= \sum_{t,q,r} |P_{T_i, V_i, Q_i, R_i}(t, 0, q, r) - P_{T_i, V_i, Q_i, R_i}(t, 1, q, r)| \\&= \sum_{t,q,r} |P(t, q)P(0, r) - P(t+1, q)P(1, r)| \\&\stackrel{(*)}{\leq} \frac{1}{2} \sum_{t,q,r} P(q) |P(0, r) - P(1, r)| + P(r) |P(t, q) - P(t+1, q)| \\&= \frac{1}{2} \sum_{t,r} |P(0, r) - P(1, r)| + \frac{1}{2} \sum_{t,q} |P(t, q) - P(t+1, q)| \\&= 2K_n,\end{aligned}$$

Identity for (*): For any a, b, c, d :

$$ab - cd = \frac{(a+c)(b-d) + (b+d)(a-c)}{2}$$

Polarization of \hat{K}_n (memory)

- ▶ Follows steps of memoryless case
- ▶ Requires additional inequalities
 - ▶ **Inequality I:** For states $s_0, s_N, s_{2N} \in \mathcal{S}$,

$$\begin{aligned} P(s_0, s_N, s_{2N}) &= \frac{P(s_0, s_N) \cdot P(s_N, s_{2N})}{P(s_N)} \\ &\leq \psi \cdot P(s_0, s_N) \cdot P(s_N, s_{2N}) \end{aligned}$$

where

$$\psi = \max_s \frac{1}{\pi(s)}$$

- ▶ **Inequality II:** For $f, g \geq 0$,

$$\sum_{s_N} f(s_N)g(s_N) \leq \sum_{s_N} f(s_N) \sum_{s'_N} g(s'_N)$$

Connections

Extreme Values

$$H \approx 0 \Leftrightarrow Z \approx 0 \Leftrightarrow K \approx 1$$

$$H \approx 1 \Leftrightarrow Z \approx 1 \Leftrightarrow K \approx 0$$

also for $(\hat{\cdot})$ processes

Ordering

$$\hat{H}_n \leq H_n$$

$$\hat{Z}_n \leq Z_n$$

$$\hat{K}_n \geq K_n$$

All six processes $(H_n, \hat{H}_n, Z_n, \hat{Z}_n, K_n, \hat{K}_n)$ polarize **fast** both to 0 and 1 with any $\beta < 1/2$

Summary

- ▶ A general framework for memory:

$$P(x_i, y_i, s_i | s_{i-1})$$

- ▶ Memory allowed in both source and channel
- ▶ State sequence S_i
 - ▶ Hidden
 - ▶ Stationary
 - ▶ Finite state Markov
 - ▶ Aperiodic and irreducible
- ▶ Achieve rate $I_*(X; Y)$ through polar codes
- ▶ No change to polarization exponent ($\beta < 1/2$)